

# Constraint Programming (CP) and IBM CP Optimizer

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Speaker: Xinbo Wang

# Outline

- **A glimpse of Constraint Programming (CP)**
- **A glimpse of Integer Mathematical Programming (MP)**
- **Comparison of CP and MP**
- **A brief introduction of IBM CP optimizer**

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## What is constraint programming?

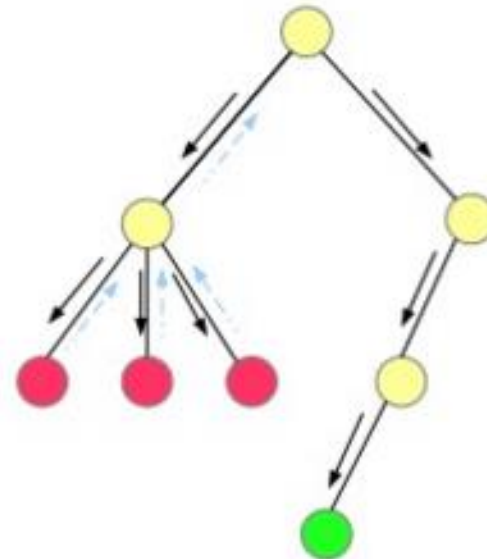
- **CP is an optimization technology which is complementary to Mathematical programming (e.g. ILP) taking a different approach to optimization, but sharing similarities.**
- **It is a relatively new technology developed in the computer science and artificial intelligence communities.**
- **It has found an important role in scheduling and highly combinational problems (for ours).**

## Applications

- Job shop scheduling
- Assembly line smoothing and balancing
- **Cellular frequency assignment**
- Airline crew rostering Nurse scheduling
- Shift planning
- Maintenance planning
- and scheduling
- Airport gate allocation and stand planning
- Production **scheduling**
- Transport **scheduling**
- Warehouse management
- Course timetabling

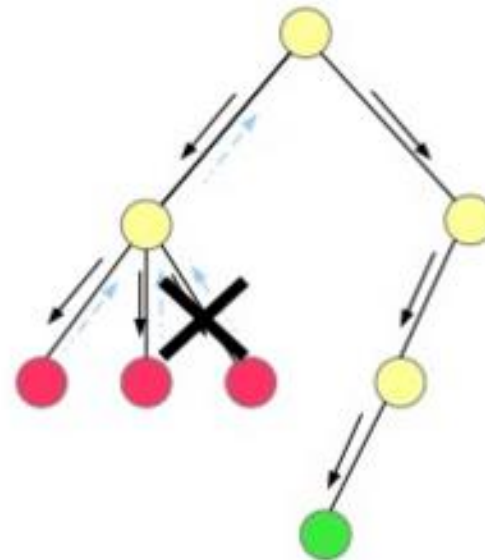
# How Constraint Programming Works?

- CP is a *constructive* approach
- Values are assigned to variables one at a time to extend a partial solution to a complete solution
- At a point, it may be useless to further extend a partial solution as at least one constraint is already violated by the partial solution
  - The solver *backtracks* and tries a different value for a previously assigned variable
  - All possible assignments of values to variables can be examined in this way



# How Constraint Programming Works?

- In CP, the basic search behaviour is tree search
- Including search space reduction via *domain filtering*
- Domain filtering
  - Before each value-variable assignment, *domain filtering* occurs
  - Each value of a variable which cannot be used in a solution (given the current partial assignment) can be removed
  - Each constraint type has a *specialized* algorithm which filters domains



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## What is Integer Linear Programming?

- An **integer programming** problem is a mathematical optimization or feasibility program in which some or all of the variables are restricted to be integers.
- The objective function and the constraints (other than the integer constraints) are linear.
- **Mixed integer linear programming** (MILP) involves problems in which only some of the variables are constrained to be integers, while other variables are allowed to be non-integers.
- **Zero-one linear programming** involves problems in which the variables are restricted to be either 0 or 1. Note that any bounded integer variable can be expressed as a combination of binary variables. For example, given an integer variable,  $x$ , the variable can be expressed :

$$x = x_1 + 2x_2 + 4x_3 + \dots + 2^{\lfloor \log_2 U \rfloor} x_{\lfloor \log_2 U \rfloor + 1}.$$

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## Comparison

- **CP works with the same concepts as mathematical programming: decision variables, objective function, and constraints.**
- **CP only discrete decision variables (integer or Boolean) vs MP discrete and continuous decision variables.**
- **CP logical constraints and arithmetic expressions (modulo, integer division, etc.) vs MP models only linear constraints or quadratic convex constraints.**
- **CP no limitation on the arithmetic constraints that can be set on decision variables vs MP specific to a class of problems whose solution space satisfies certain mathematical properties.**
- **Each optimization engine uses different techniques and algorithms to find feasible solutions and optimize them.**

## A Tabular View

### *Constraint programming vs. mathematical programming*

Feature	MP	CP
Relaxation	Yes	No
GAP measure	Yes	No
Optimality proof	Yes	Yes
Modeling limitations	Quadratic problems are limited to PSD (Positive Semi Definite) problems and Second Order Cone Programming (SOCP) problems	Discrete problems
Specialized constraints	No	Yes
Logical constraints	Yes	Yes
Theoretical grounds	Algebra	Graph theory and algorithmic
Modeler support	Yes	Yes
Model and run	Yes	Yes

## Benefits of constraint programming

- **Solve time tabling problems and sequencing problems.**
- **An alternative to mathematical programming for allocation problems that have a slow convergence.**
- **Constraint programming has native support for:**
  - ✓ **Nonlinear costs or constraints**
  - ✓ **Logical constraints and statements**
  - ✓ **Constraints on and between interval variables**
  - ✓ **Compatibility or incompatibility constraints**
  - ✓ **More useful features**

## Expressions and Constraints

- **Arithmetic constraints**

- ✓  $x + y, x - y, x * y, x / y, x \text{ div } y, x \% y$
- ✓  $\min, \max, \text{abs}, \log, \exp$  etc.
- ✓ Piecewise linear functions

- **Relational constraints**

- ✓  $x == y, x != y, x <= y, x < y, lb <= x <= ub$

- **Logical constraints**

- ✓  $!c, c || d, c \&\& d,$
- ✓  $c \Rightarrow d, c \Rightarrow d \text{ else } e$
- ✓  $c$  and  $d$  are relational or conditional constraints

## Expressions and Constraints

- **Reification**

- ✓ Relational or logical constraints can be used in a value context, where they evaluate to 0 or 1

- **Examples**

- ✓ Arithmetic:  $\max(0, \text{abs}(\text{load}[i] - \text{cap}))$
- ✓ Relational:  $\text{wid} * \text{hei} * \text{depth} * \text{density} \leq \text{maxLoad}$
- ✓ Logical:  $\text{end}[i] \leq \text{start}[j] \vee \text{start}[j] \leq \text{end}[i]$
- ✓ Reification:  $\text{spill} == (\text{load}[i] > \text{cap})$

## Expressions and Constraints

- **Count expression**

- ✓ `count(dvar int[] x, int c)`
- ✓ Evaluates the number of variables in `x` with value `c`
- ✓ e.g. Count the number of nurses allocated to ward 5
  - `count(wardAllocation, 5) >= 3`

- **Element expression**

- ✓ `(int[] a)[dvar int x]` OR `(dvar int[] a)[dvar int x]`
- ✓ Evaluates to the `x`th member of `a`
- ✓ e.g. `travel == 2 * distFromPittsburgh[holidayTown]`
- ✓ `travel` and `holidayTown` are variables



## Expressions and Constraints

- **All Different**
  - ✓ `allDifferent(dvar int[] x)`
  - ✓ All variables in `x` must take different values
  - ✓ e.g. The rank (visit priority) of each city is different
    - `allDifferent(rankOfVisit)`
- **Allowed / Forbidden assignments**
  - ✓ `allowedAssignments({<a,b,c>} A, dvar int[3] x)`
  - ✓ The assignments to `x` must fit with a tuple of `A`
  - ✓ `forbiddenAssignments` is the negation of this

## Expressions and Constraints

- **Bin packing constraint**
  - ✓ `pack(dvar int[m] ld, dvar int[n] x, int[n] sz, dvar int c)`
  - ✓ `ld[i] == sum(j) (x[j] == i) * sz[j]`
    - `c` is the number of containers used
- **Inverse constraint**
  - ✓ `inverse(dvar int[n] x, dvar int[n] y)`
  - ✓ `x[i] == j <=> y[j] == i` --- link primary and dual models
- **Lexicographic ordering constraint**
  - ✓ `lex(dvar int[n] x, dvar int[n] y)` --- break symmetries

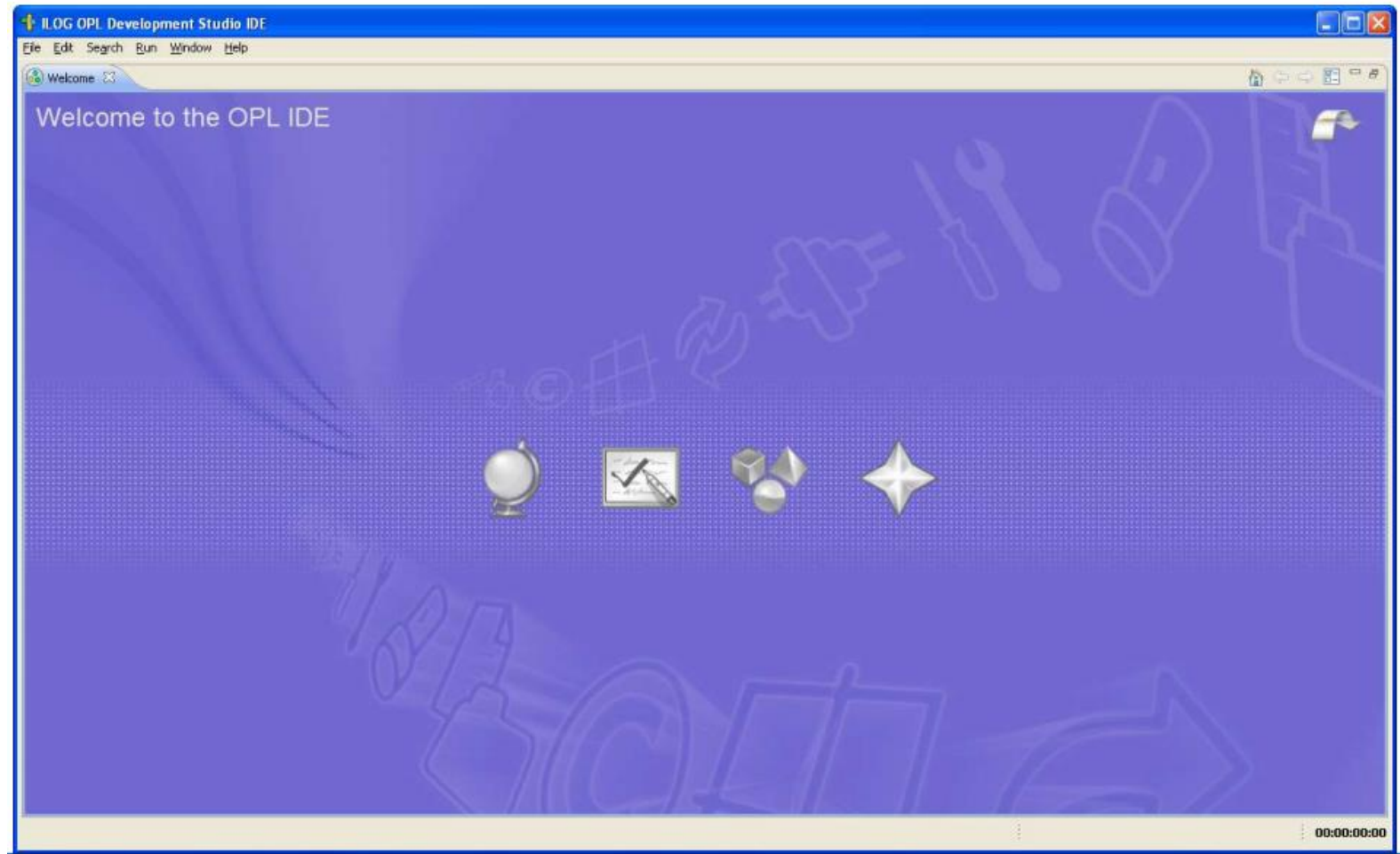
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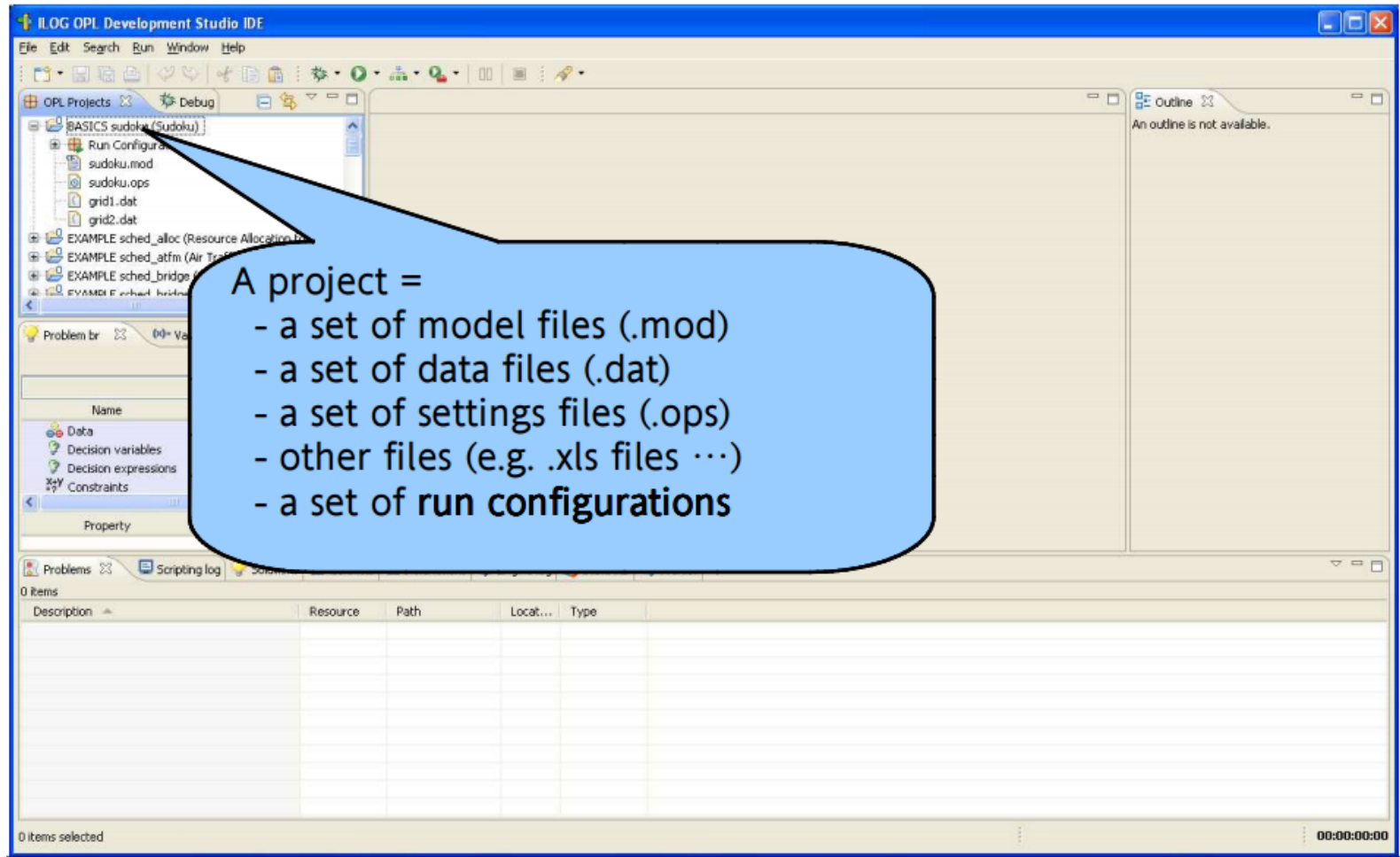
## What is CP Optimizer

- **A Constraint Programming engine with an emphasis on modelling and automatic search**
- **Available as a toolkit in C++ , Java, .NET**
  - ✓ C++ is the native language and allows more possibilities, like writing incremental custom constraints, and fully controlling the search process
- **Available as an engine inside ILOG OPL IDE**
  - ✓ ILOG: name of the company (acquired by IBM)
  - ✓ OPL: optimization programming language
  - ✓ Higher level modelling and data manipulation

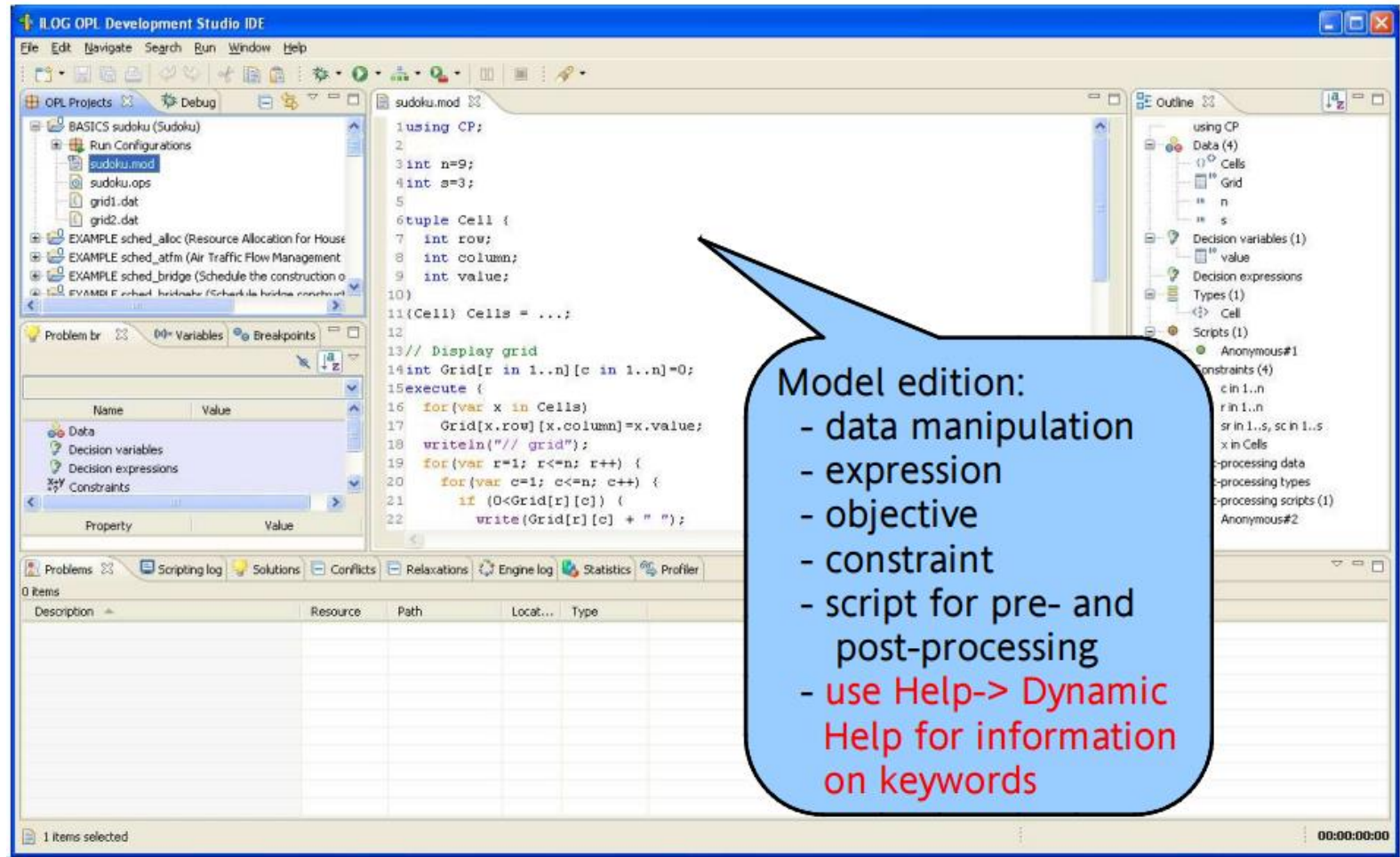
# ILOG OPL IDE



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# ILOG OPL IDE



# ILOG OPL IDE

The screenshot shows the ILOG OPL Development Studio IDE. The main window displays a Sudoku puzzle model. The code editor shows a list of cells with their coordinates and values. A data grid is overlaid on the code editor, showing the Sudoku puzzle grid. The project explorer on the left shows the project structure, including a file named 'grid1.dat'. A blue callout box points to 'grid1.dat' and contains the following text:

Data file:

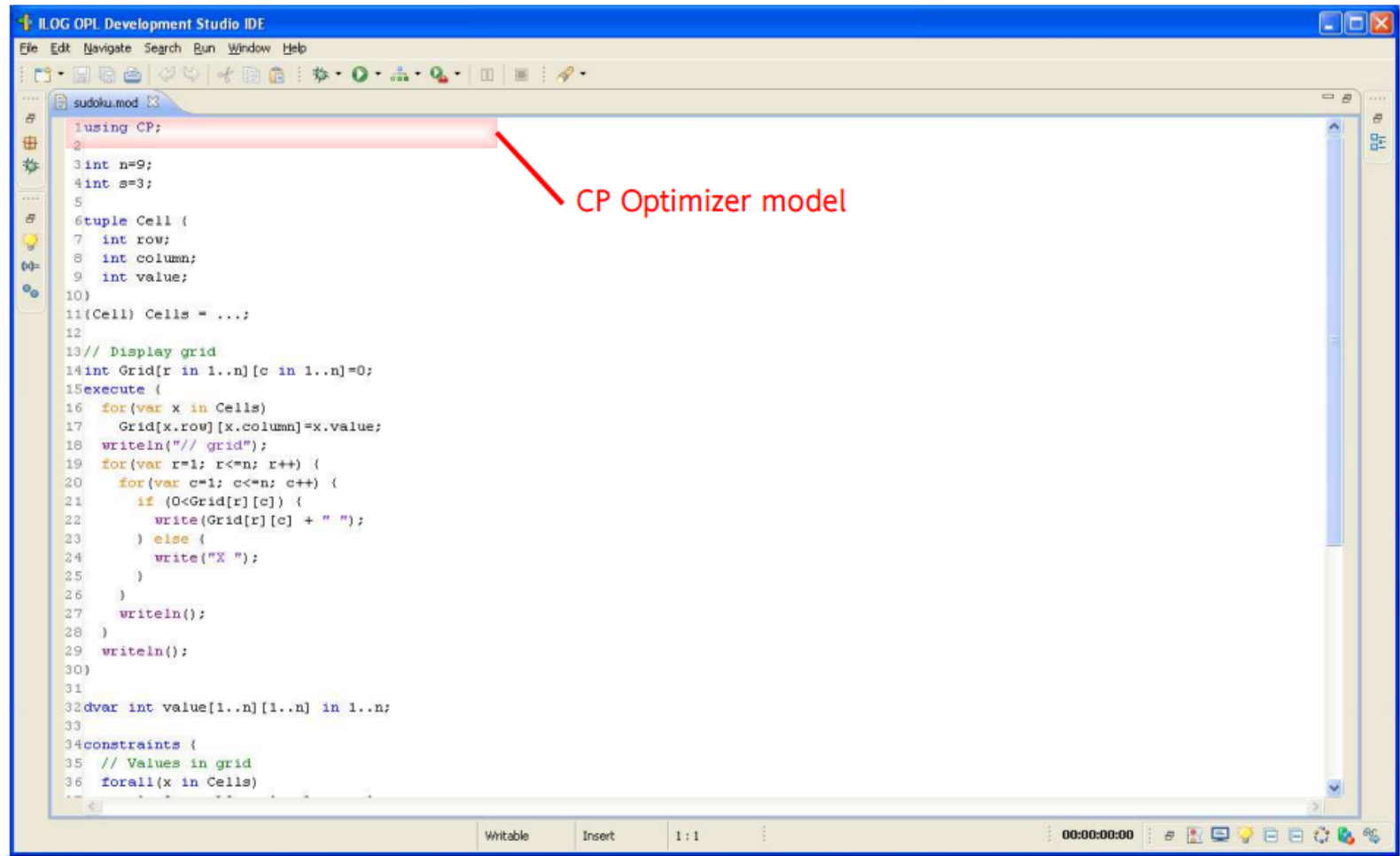
- explicit data
- connection to:
  - \* databases
  - \* Excel worksheets



# Overview of an CP Model using OPL

- **Top**
  - ✓ Data manipulation and pre-processing
    - declarative (expressions) and/or imperative (script)
  - ✓ Variable declarations
- **Middle**
  - ✓ Declarative model
    - objective (optional) and constraints
- **Bottom**
  - ✓ Post-processing of solutions
  - ✓ Declarative (expressions) and/or imperative (script)

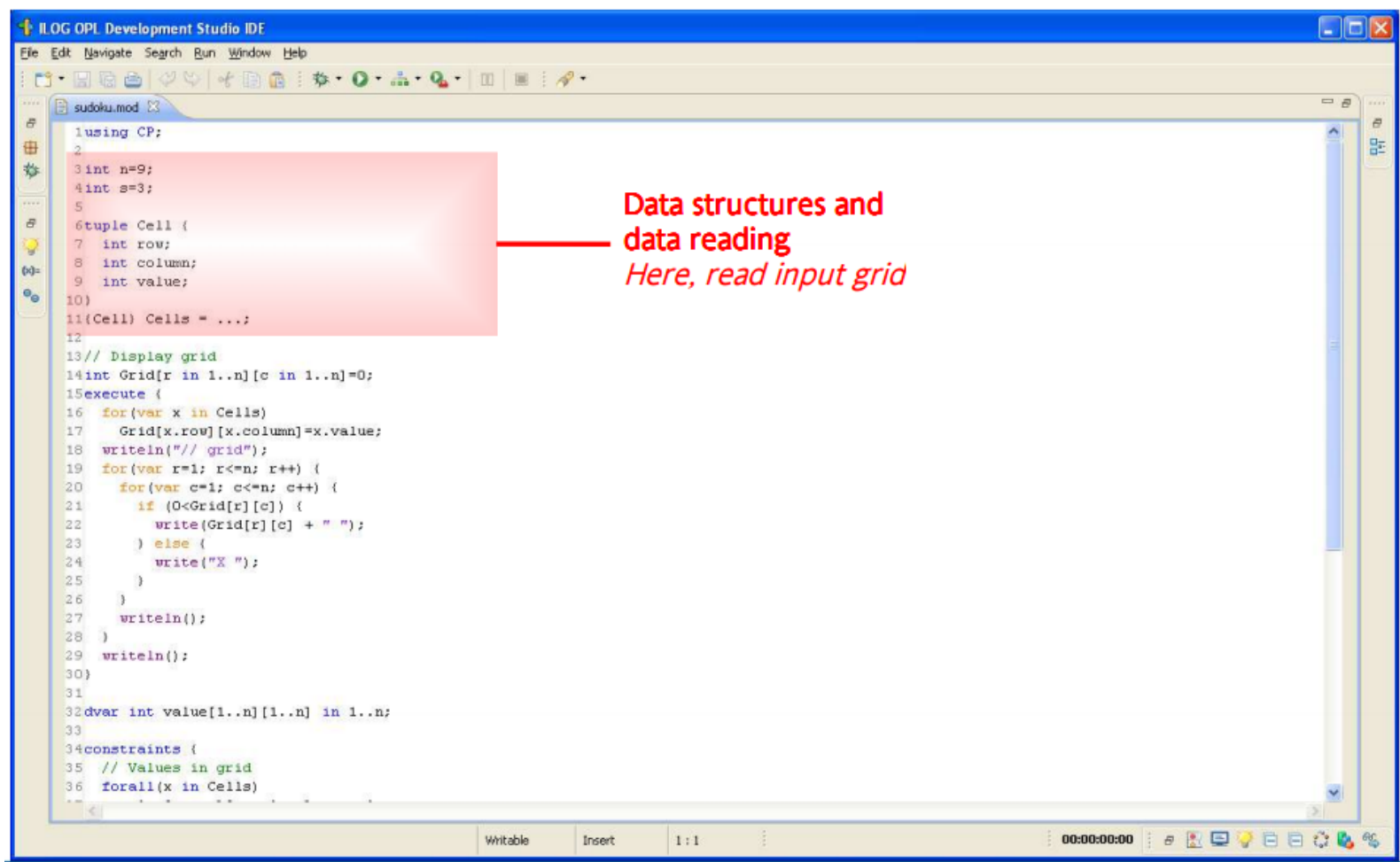
# An CP Model



The screenshot shows the ILOG OPL Development Studio interface. The main window displays a file named 'sudoku.mod'. A red arrow points from the text 'CP Optimizer model' to the first line of the code, 'using CP;'. The code is as follows:

```
1 using CP;
2
3 int n=9;
4 int s=3;
5
6 tuple Cell {
7   int row;
8   int column;
9   int value;
10}
11(Cell) Cells = ...;
12
13// Display grid
14int Grid[r in 1..n][c in 1..n]=0;
15execute {
16   for(var x in Cells)
17     Grid[x.row][x.column]=x.value;
18   writeln("// grid");
19   for(var r=1; r<=n; r++) {
20     for(var c=1; c<=n; c++) {
21       if (0<Grid[r][c]) {
22         write(Grid[r][c] + " ");
23       } else {
24         write("X ");
25       }
26     }
27     writeln();
28   }
29   writeln();
30}
31
32dvar int value[1..n][1..n] in 1..n;
33
34constraints {
35   // Values in grid
36   forall(x in Cells)
```

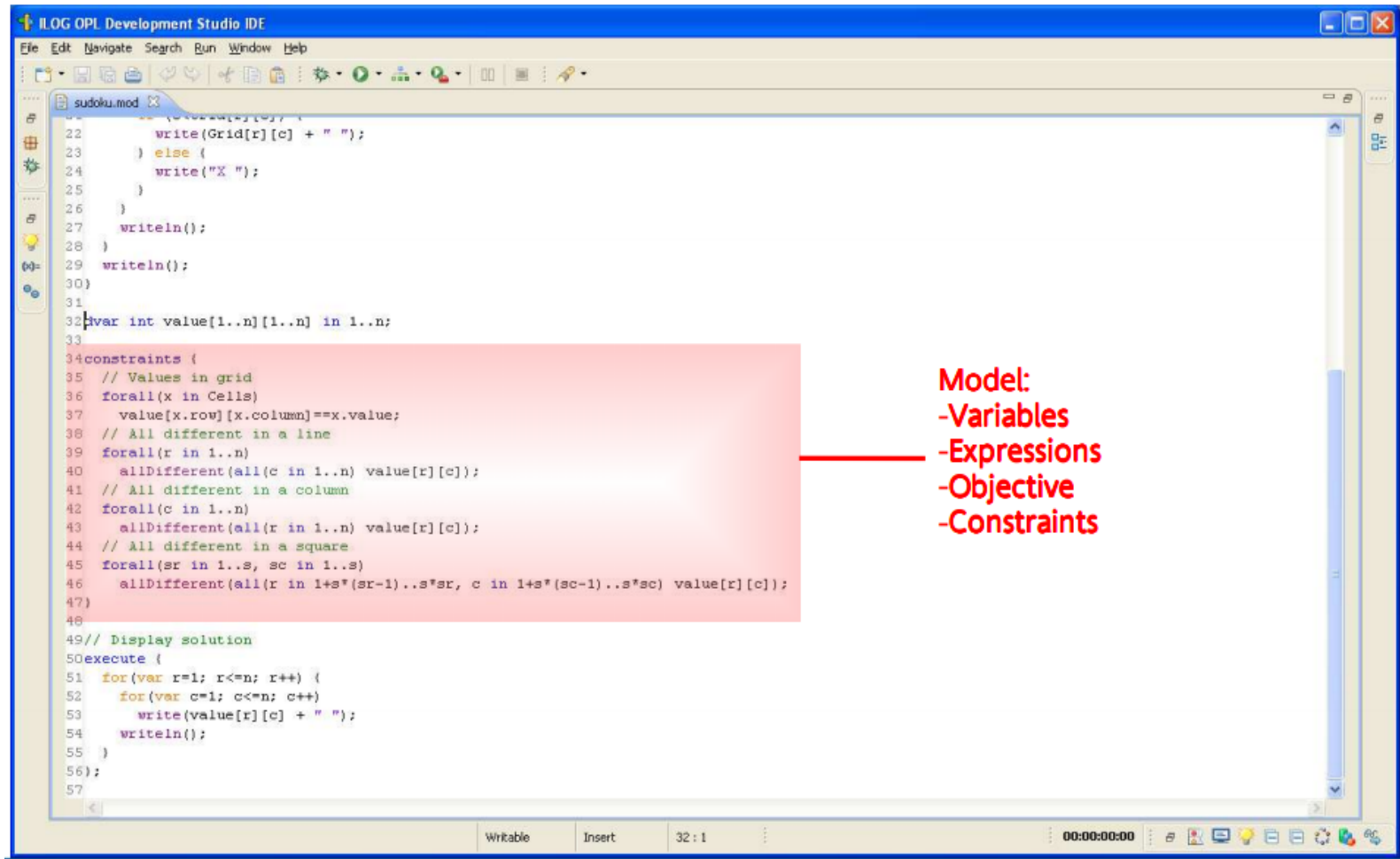
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Data structures and  
data reading  
*Here, read input grid*

# An CP Model



```

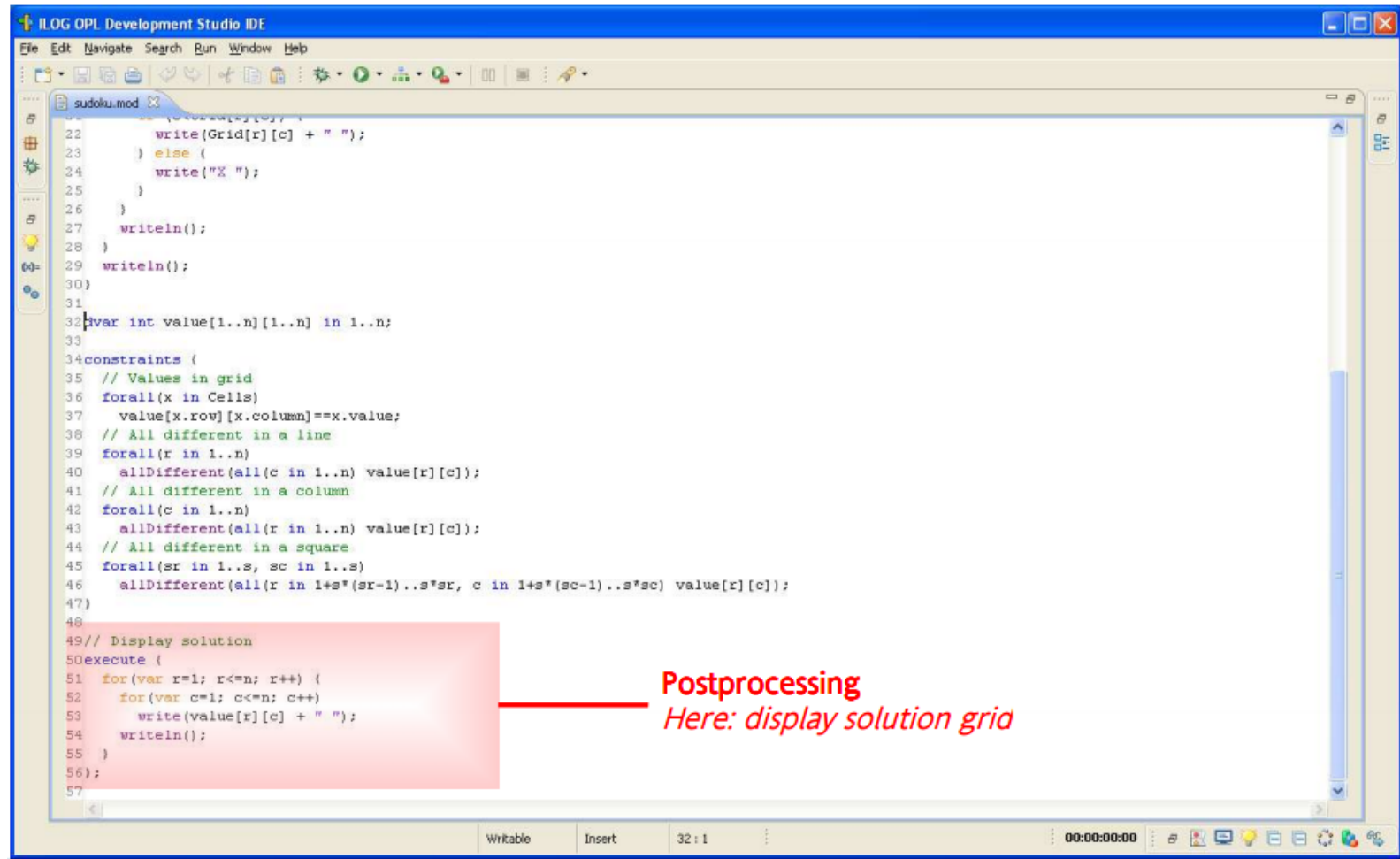
1  {0..n-1}{0..n-1}
22  write(Grid[r][c] + " ");
23  } else {
24    write("X ");
25  }
26  }
27  writeln();
28  }
29  writeln();
30  }
31
32  var int value[1..n][1..n] in 1..n;
33
34  constraints {
35    // Values in grid
36    forall(x in Cells)
37      value[x.row][x.column]==x.value;
38    // All different in a line
39    forall(r in 1..n)
40      allDifferent(all(c in 1..n) value[r][c]);
41    // All different in a column
42    forall(c in 1..n)
43      allDifferent(all(r in 1..n) value[r][c]);
44    // All different in a square
45    forall(sr in 1..s, sc in 1..s)
46      allDifferent(all(r in 1+s*(sr-1)..s*sr, c in 1+s*(sc-1)..s*sc) value[r][c]);
47  }
48
49  // Display solution
50  execute {
51    for(var r=1; r<=n; r++) {
52      for(var c=1; c<=n; c++)
53        write(value[r][c] + " ");
54      writeln();
55    }
56  };
57

```

**Model:**

- Variables
- Expressions
- Objective
- Constraints

# An CP Model



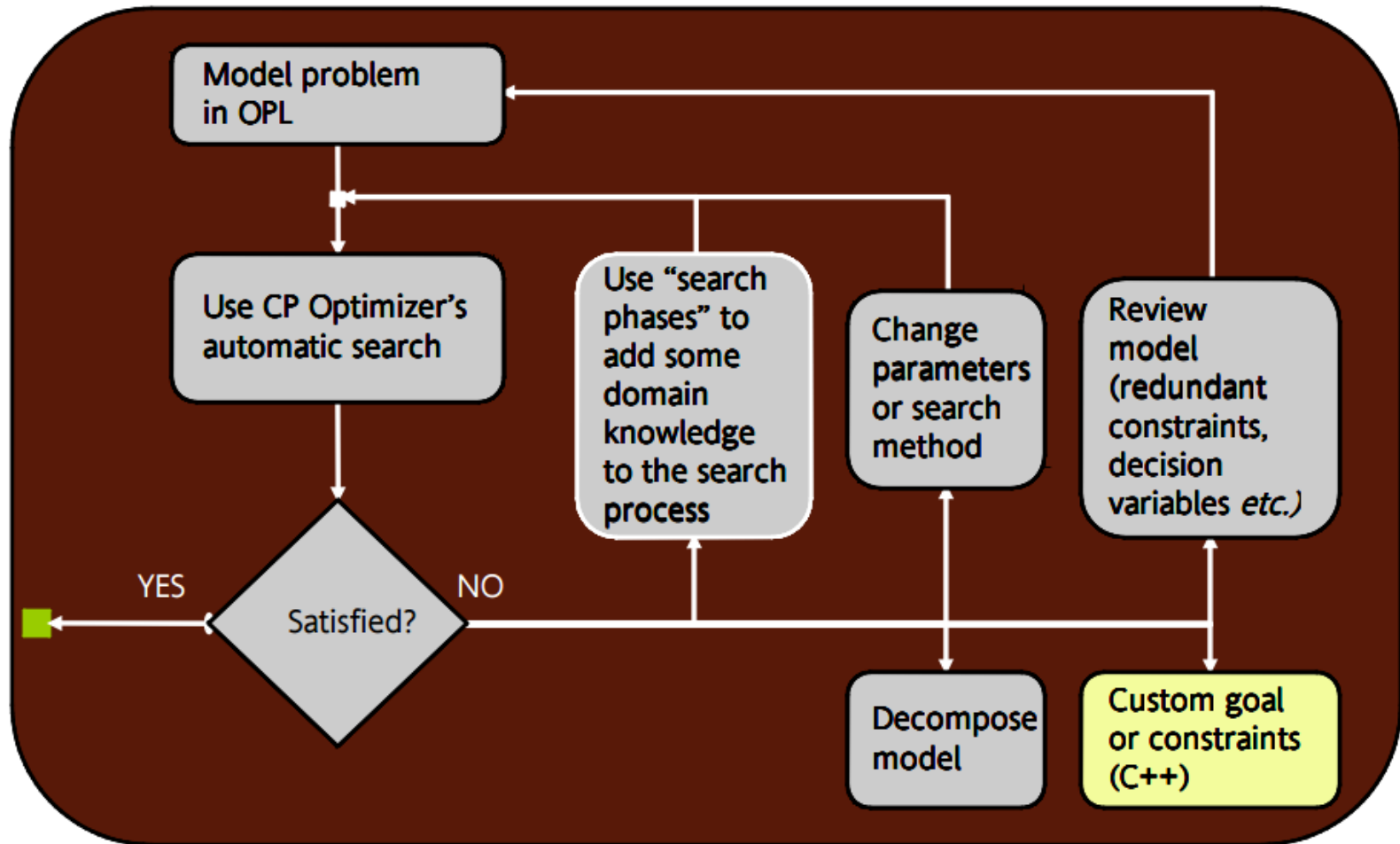
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54     writeln();
55   }
56 }
57
```

Postprocessing  
Here: display solution grid

## Search in CP Optimizer

- **Automatic search is emphasized**
  - ✓ Simpler, more maintainable, benefit from upgrades
- **Search Phases**
  - ✓ What group of variables to assign first
  - ✓ (optionally) define instantiation strategy
- **Parameters**
  - ✓ Inference levels and search control parameters
- **Problem still hard?**
  - ✓ Improve model
  - ✓ Simplify or relax specification
  - ✓ Decompose: CPLEX often useful here

## Typical Use of CP Optimizer



Thank You !

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