

Applying Hadoop Cloud Computing Technique to Optimal Design of Optical Networks

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Group meeting, UC. Davis, 2/5/2016

Outline



- Background
- Hadoop cloud computing system
- Optimization for optical networks with Hadoop system
- Simulation and performance analyses
- Conclusions

Background



Bin-packing problem

A classical NP-hard problem that targets to find a minimal number of bins to pack all the (irregular) items



Sequence 1





Sequence 2



Background (cont's)



- Bin-packing problems in optical networks
 - Energy minimization problem, shared backup path protection (SBPP) planning problem, adaptive FEC assignment problem, etc.
 - → For a traffic demand, in addition to its size (i.e., traffic volume), we also need to consider its shape (i.e., the route of the demand) when packing (i.e., serving) them
 - → For these research problems like serving network traffic demands, the simple largest-smallest ordering strategy cannot guarantee good performance



Background (cont's)



- Impact of demand sequence
 - → To achieve good performance, the most straightforward strategy is to randomly shuffle the demand sequences
 - Then implement an efficient heuristic algorithm for each of them and finally choose the one with the best performance



However, it becomes computationally intensive and time-consuming if a huge number of shuffuled demand sequences are evaluated

[1] G. Shen, Y. Li, and L. Peng, "Almost-optimal design for optical networks with Hadoop cloud computing: ten ordinary desktops solve 500-node, 1000-link, and 4000-request RWA problem within three hours (Invited)," in Proc. *ICTON 2013*.



- Features and advantages
 - → Open source software framework
 - → Employ many low-end ordinary machines
 - Provide high parallel computation capacity
 - → Cost effective
- Key functional modules
 - Hadoop Distributed File System (HDFS)
 - → Map/Reduce process



- Hadoop Distributed File System (HDFS)
 - A highly fault-tolerant distributed file system designed for commodity hardware
 - A HDFS instance typically includes a Namenode and many Datanodes
 - Namenode is a coordinator and Datanodes provide storage for HDFS





- Map step
 - → Receive problem data
 - \rightarrow Divide the problem into smaller sub-problems
 - Distribute them to Slave nodes which will solve the subproblems and feedback their obtained results to the Master node

Reduce step

→ Master node is employed to collect the answers to all the sub-problems and combine them in some way to form a final solution



- Hardware
 - → Lenovo YangTian T4900d (Dual cores 2.9-GHz CPU)
- Software
 - → Linux (Ubuntu 12.04 LTS)
 - → Hadoop-1.0.3
 - → JDK: Sun-java6-jdk
 - \rightarrow SSH











- Seven ordinary desktops Hadoop system
 - One master and seven slaves interconnected by a switch
 - → All the computers have 2.9-GHz dual cores with 2G memory
 - Two reduces on each slave node





- Energy minimization problem [2]
 - → A "Follow the Sun, Follow the Wind" strategy for the IP over WDM network
 - An efficient heuristic algorithm to minimize the total nonrenewable energy consumption of the network

→ Shuffling strategy

- → Shuffle the node pairs in a list to change their order, and recalculate non-renewable energy consumption
- → Compare the results and select the least non-renewable energy as our final result

An example with shuffling strategy:

- (a) Traffic matrix with serving order 1
- (b) Established virtual topology for serving order 1
- (c) Traffic matrix with serving order 2
- (d) Established virtual topology for serving order 2









[2] G. Shen, Y. Lui, and S. K. Bose, ""Follow the Sun, Follow the Wind" lightpath virtual topology reconfiguration in IP over WDM network," *IEEE/OSA Journal of Lightwave Technology*, vol. 32, no. 11, pp. 2094-2105, Jun. 2014.



The diagram of Map/Reduce process





Map/Reduce process

→ Map:

- → Generates a total of *N* key-value pairs (k_i, d_i) , for each of which k_i and d_i are both set to be *i*, i.e., $k_i=d_i=i$
- \rightarrow The key-value pairs are forwarded to *r* Reduce functions for parallel computing
- → Reduce:
 - \rightarrow Each of the Reduce functions has N/r key-value pairs
 - → For each key-value pair (k_i, d_i) , we randomly shuffle the demand list *M* times with value d_i as a random seed.
 - → Each of the shuffled demand sequences is then served by the heuristic "algorithm" to generate *M* results for each key-value pair (k_i, d_i)
 - → We compare these results to choose the one with the best performance for each Reduce function

Compare the best results of each Reduce function to obtain a final global optimum, i.e., "best result"

The whole process evaluates a total of $N \cdot M$ shuffled demand sequences to obtain the final best solution

Simulation conditions



- Map/Reduce process
 - →M: 10
 - → N: [100, 10,000]
 - Shuffled demand sequences: [1,000, 100,000]
- Energy minimization problem
 - Test network case: 43-link US backbone network (USNET)
 - Traffic demand between each node pair is randomly generated within the range of [10, 30] Gb/s
 - The total non-renewable energy consumption is shown for USNET on May 2th, 2014



Total non-renewable energy consumption and computation time (USNET)



- When the number of shuffled demand sequences increases from 1,000 to 10,000, the total nonrenewable energy consumption decreases 8%.
- Hadoop system can significantly shorten computation time for the same number of shuffled demand sequences
- Hadoop system is more efficient for a larger optimization problem



- SBPP network planning problem [3]
 - Distance adaptive routing and spectrum assignment (RSA) for SBPP-based elastic optical networks
 - An efficient heuristic algorithm based on spectrum window planes (SWPs) to minimize the maximal number of required frequency slots (FSs) in the entire network
 - → Shuffling strategy
 - \rightarrow Shuffle the node pair list many times
 - → Run the heuristic algorithm for each shuffled node pair list to obtain the required number of FSs, FS_i
 - → Finally compare the obtained FS_i to choose the one with the smallest number, i.e., $FS_{min} = min \{FS_i\}$.
- [3] C. Wang, G. Shen, and S. K. Bose, "Distance adaptive dynamic routing and spectrum allocation in elastic optical networks with shared backup path protection," *IEEE/OSA Journal of Lightwave Technology*, vol. 33, no. 14, pp. 2955-2964, Apr. 2015.

Simulation conditions



- SBPP network planning problem
 - → Test network case: COST239 network
 - → Frequency slots in each link: 400 FSs
 - → Bandwidth granularity of each FS: 12.5 GHz
 - → Modulation format set: BPSK, QPSK, 8-QAM
 - The bandwidth request between each node pair is uniformly randomly distributed within a range of [10, 400] Gb/s

Performance analyses



Maximal number of required FSs and computation time (COST239)



- The shuffling strategy is helpful to significantly reduce the number of required FSs
- The Hadoop system does help to reduce computation time compared to a single machine



- Adaptive FEC assignment problem [4]
 - A new adaptive FEC assignment strategy for the elastic optical network
 - → A heuristic algorithm based on slot window planes (SWPs) to maximize the total number of FSs used for actual user data transmission
 - → Shuffle strategy
 - → Shuffle the lightpath demand list for many times
 - → For each shuffled demand list, the adaptive FEC-based SWP algorithm is applied to find a solution
 - → The best solution is selected as the final solution when all the shuffled demand lists have been evaluated and compared.

[4] Y. Li, H. Dai, G. Shen, and S. K. Bose, "Adaptive FEC selection for lightpaths in elastic optical networks," in Proc. *OFC 2014*.

Simulation conditions



- Adaptive FEC assignment problem
 - Test network case: 43-link US backbone network (USNET)
 - → Frequency slots in each link: 400 FSs
 - The number of FSs required by users is uniformly randomly distributed within a range of [25, 35]

Performance analyses



 Total number of FSs for user data transmission and computation time (USNET)



- The Hadoop system can perform better than a single machine by 3% and it can also shorten computation time 20% even though it evaluates 10 times more shuffled demand sequences
- A further increase of the number of shuffled demand sequences to 1,000 can lead to 2% furtjer performance improvement, under which the computational efficiency of the Hadoop system is more significant to take only about 1 hour, while a single desktop took about 9 hours

Conclusions



- The bin-packing optimization problems in the optical network can be efficiently solved by evaluating multiple shuffled demand sequences
- To overcome the difficulty of high computational complexity, we built a Hadoop computing system to parallel evaluate multiple shuffled demand sequences
- The approach was efficient to find solutions close to optimality and significantly reduce computation times compared to a single machine



