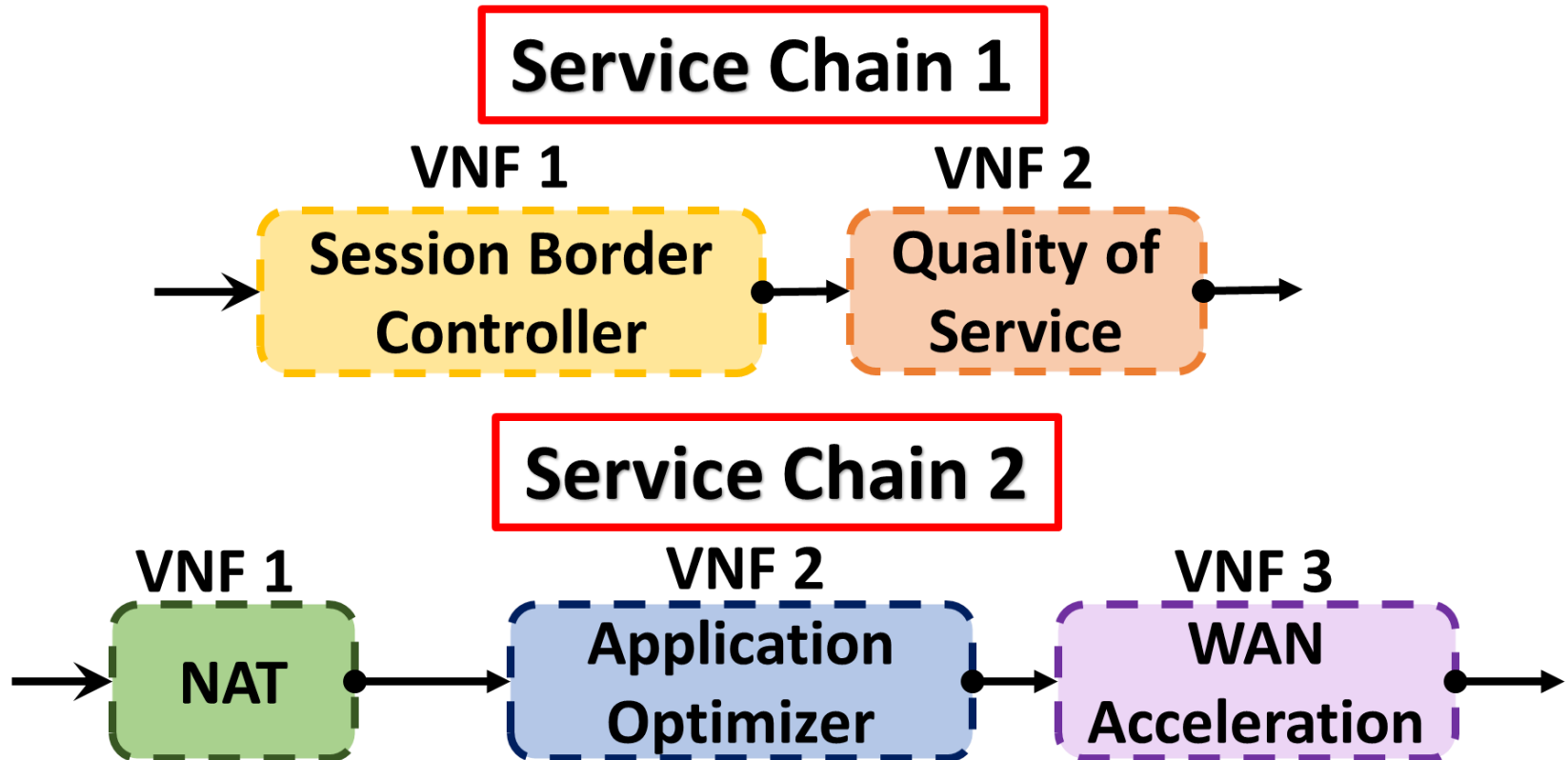


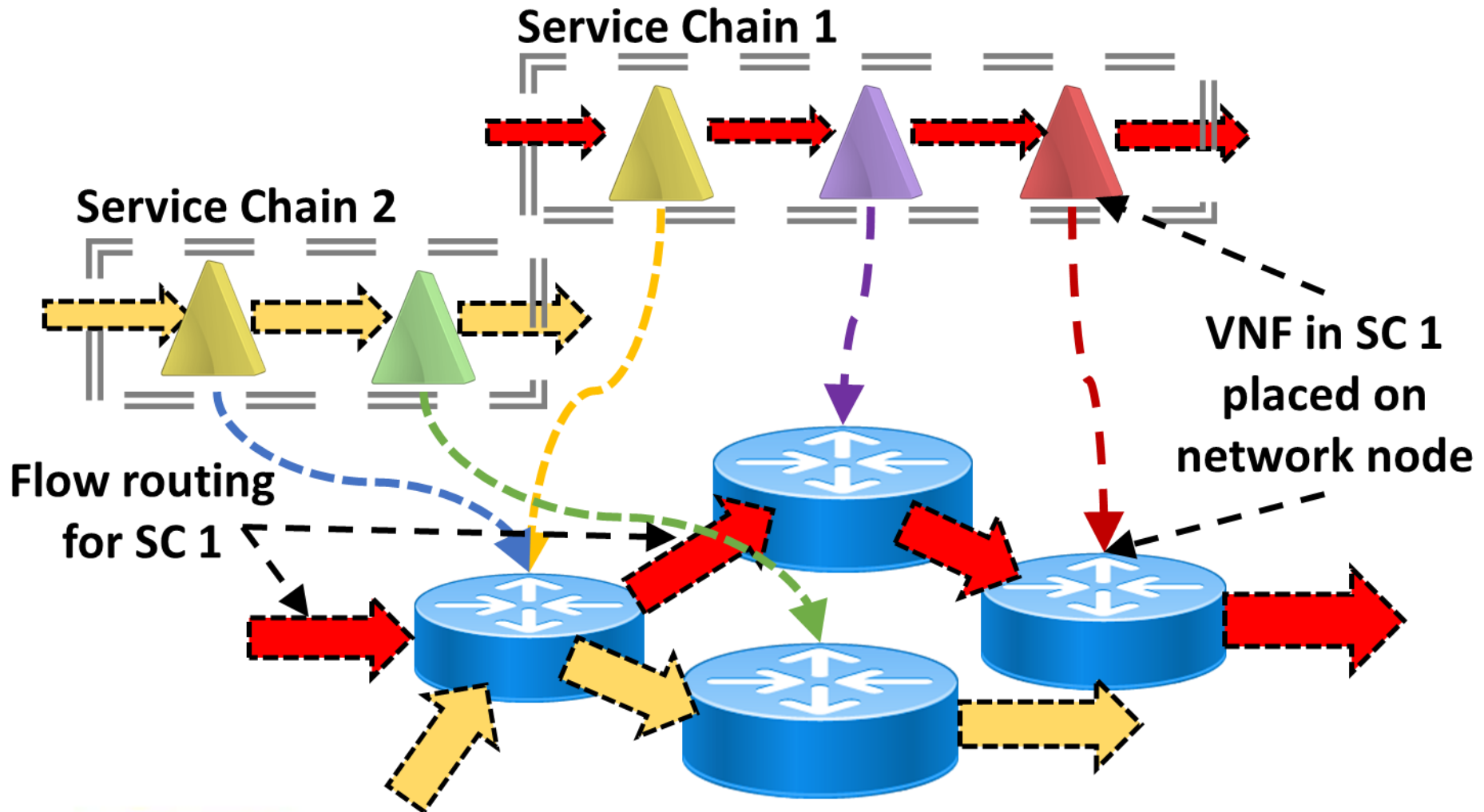
Deploying Multiple Service Chain (SC) Instances per Service Chain

BY
ABHISHEK GUPTA
FRIDAY GROUP MEETING
APRIL 21, 2017

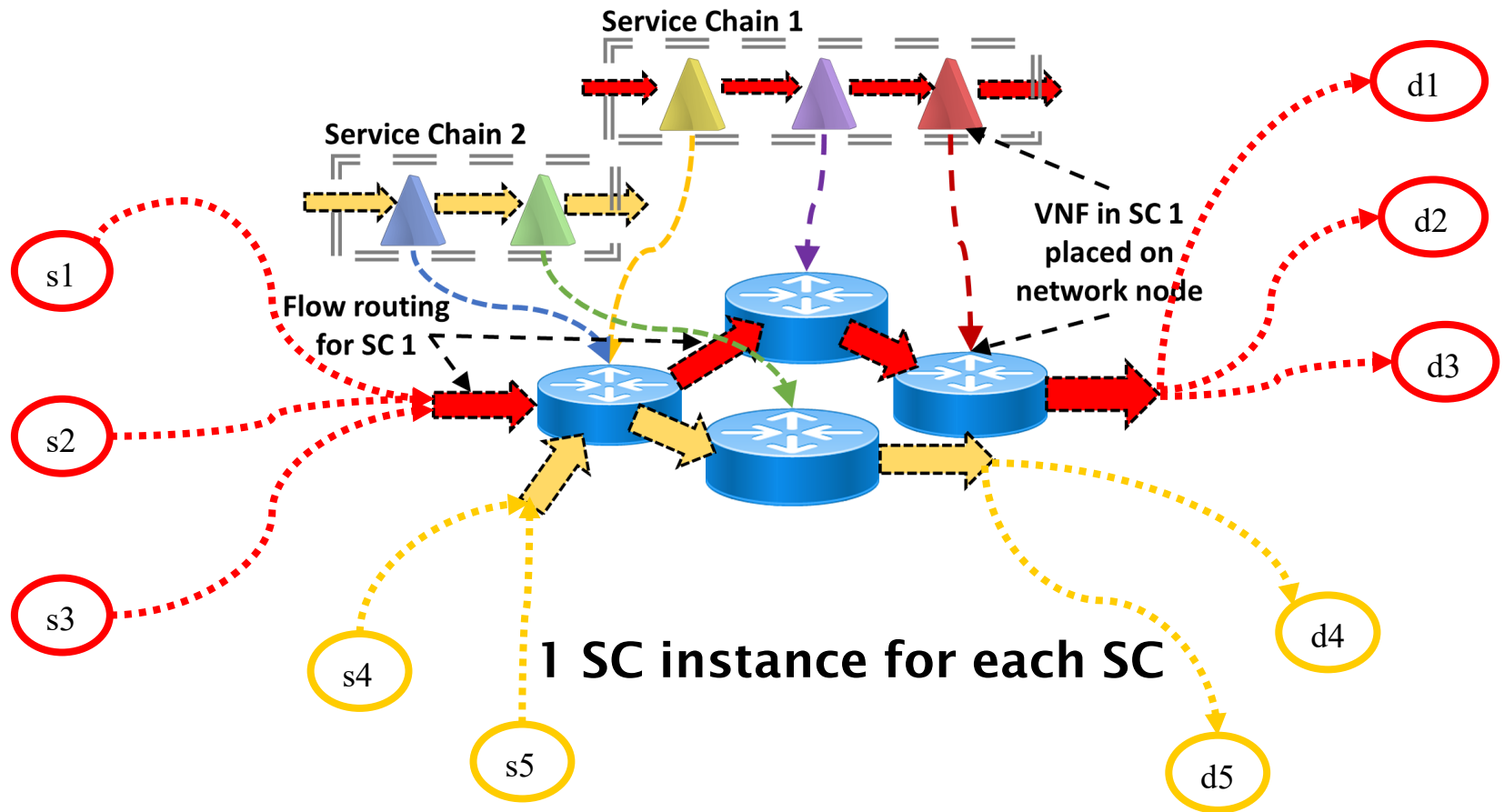
Virtual Network Function (VNF) Service Chain (SC)



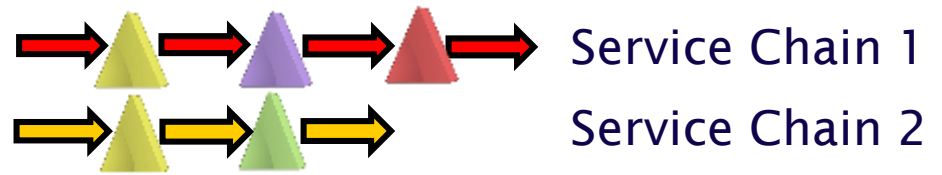
Multiple VNF SC Placement and Routing



Continued...



Single Instance Per SC



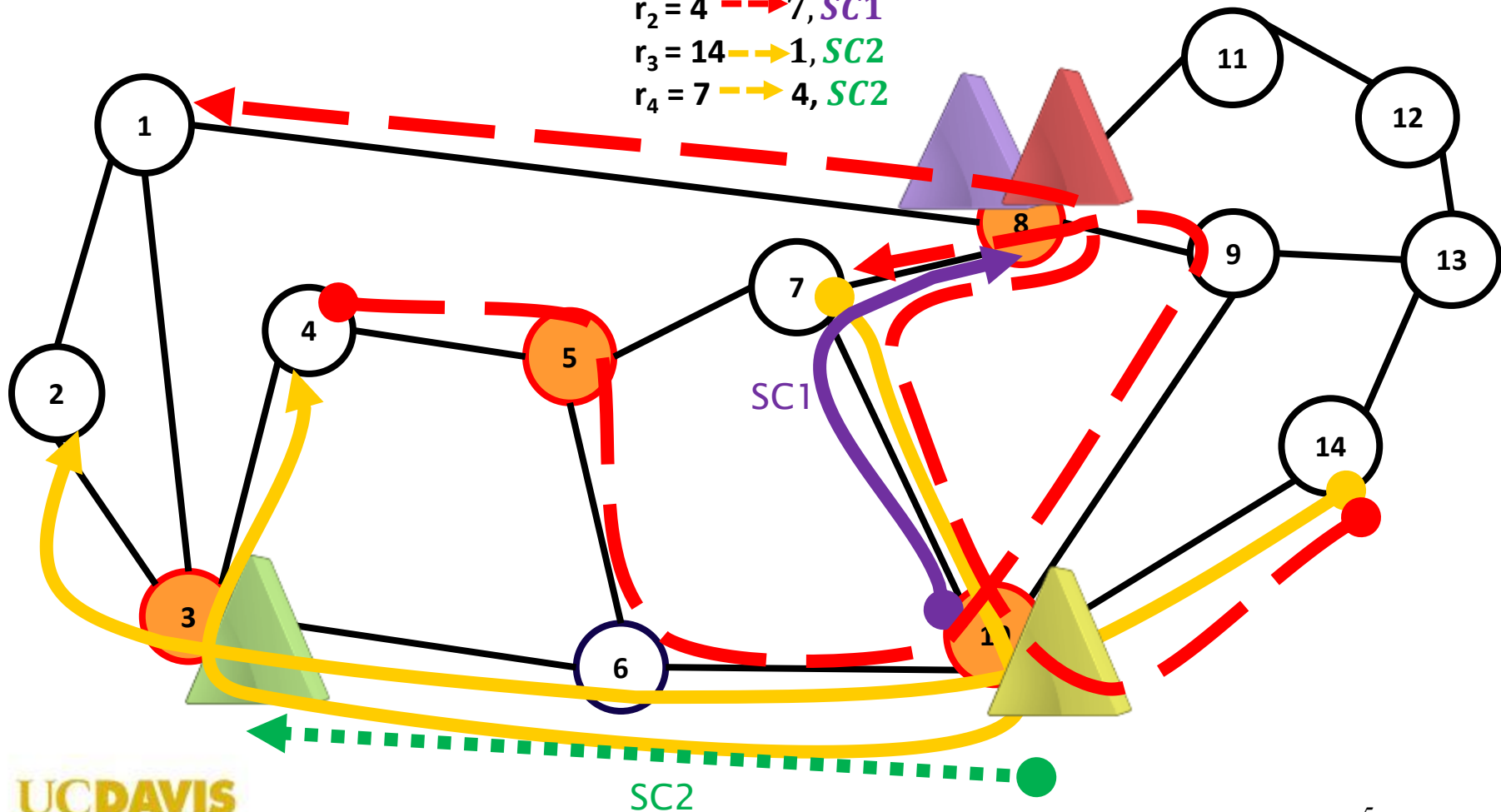
3 Demand Requests

$r_1 = 14 \rightarrow 1, SC1$

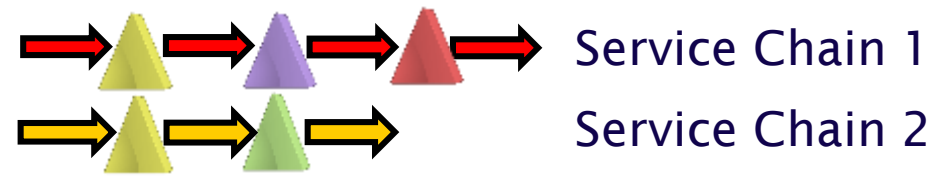
$r_2 = 4 \rightarrow 7, SC1$

$r_3 = 14 \rightarrow 1, SC2$

$r_4 = 7 \rightarrow 4, SC2$



Multiple Instances per SC



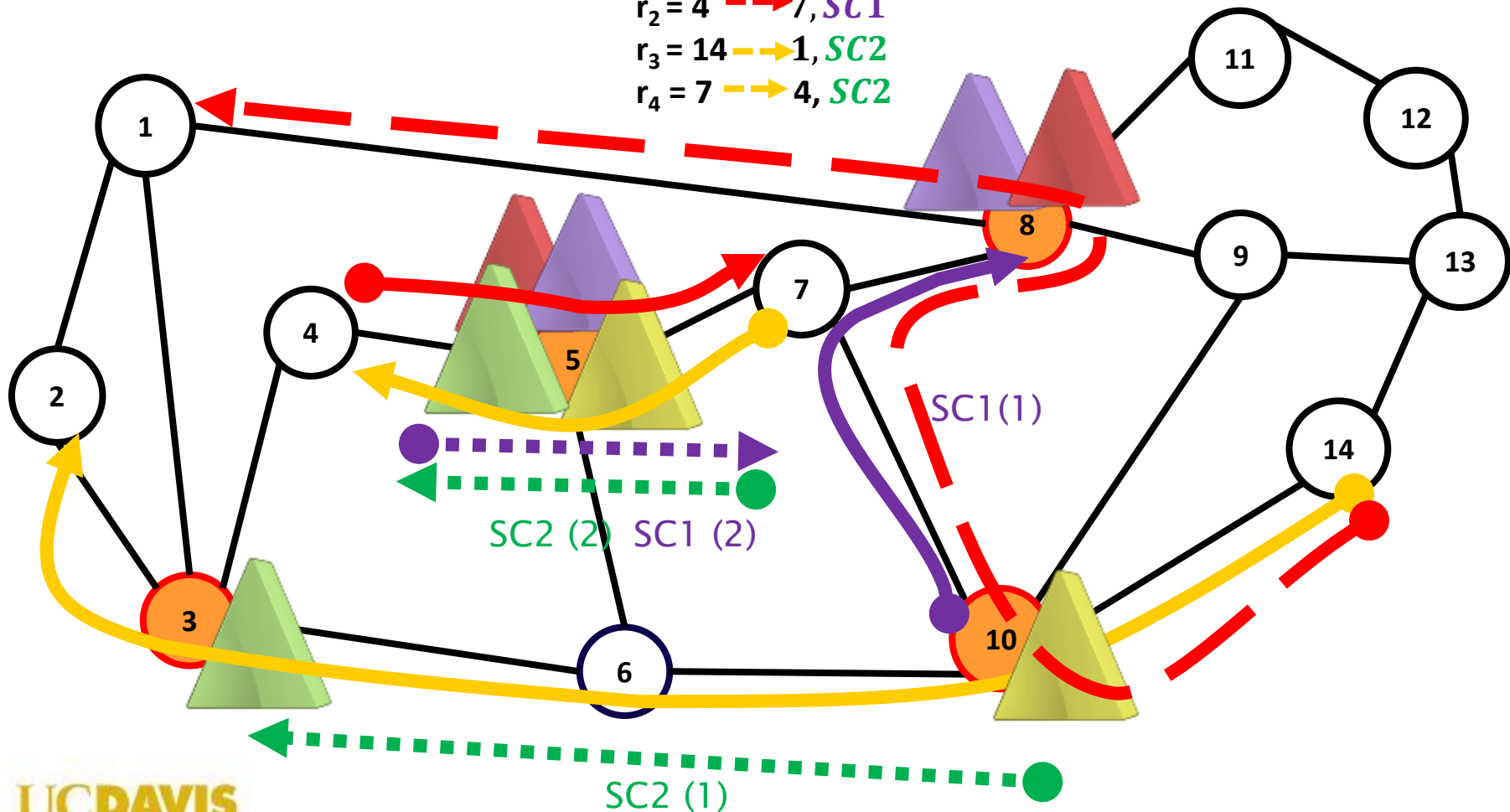
3 Demand Requests

$r_1 = 14 \rightarrow 1, SC1$

$r_2 = 4 \rightarrow 7, SC1$

$r_3 = 14 \rightarrow 1, SC2$

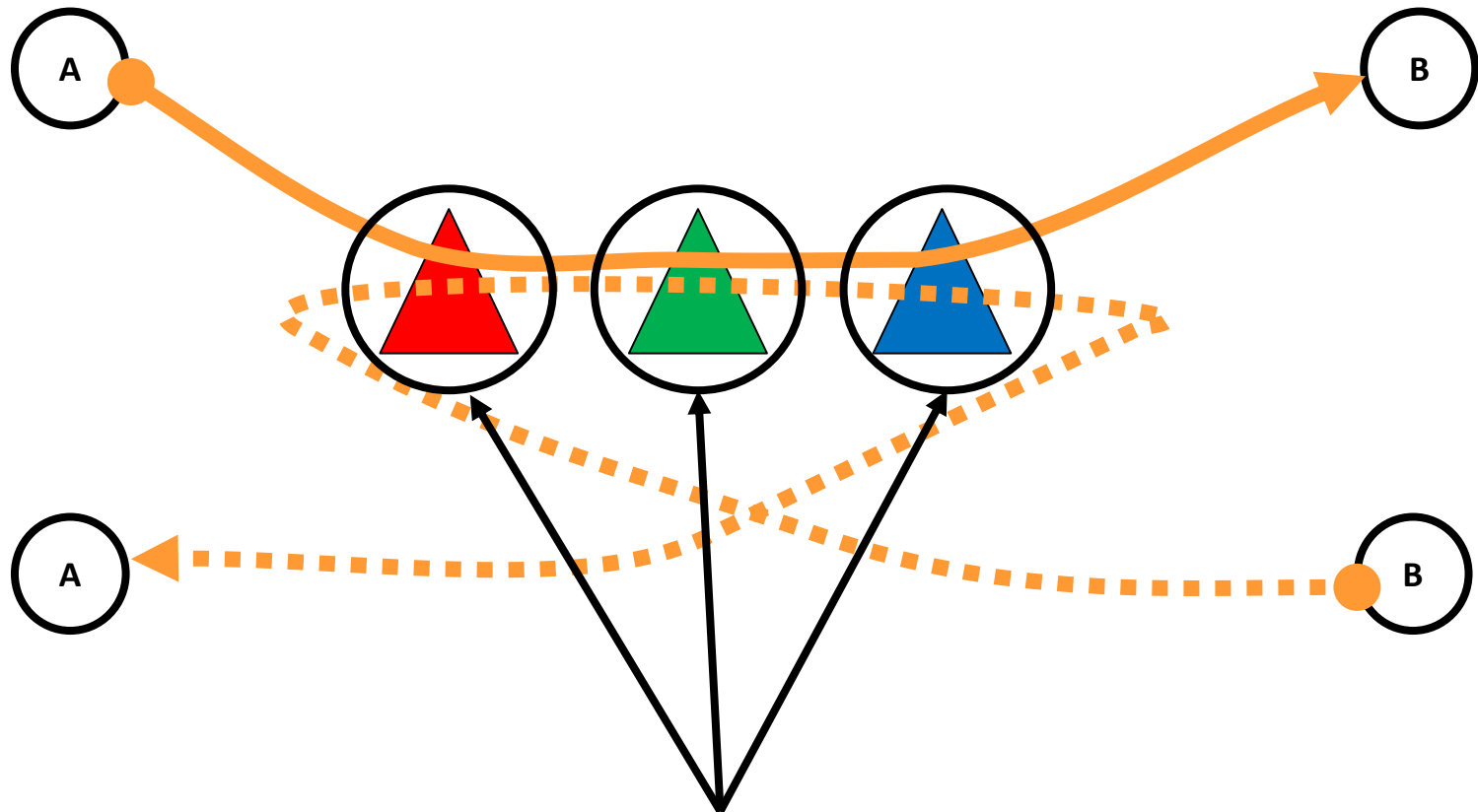
$r_4 = 7 \rightarrow 4, SC2$



Inferences and Questions

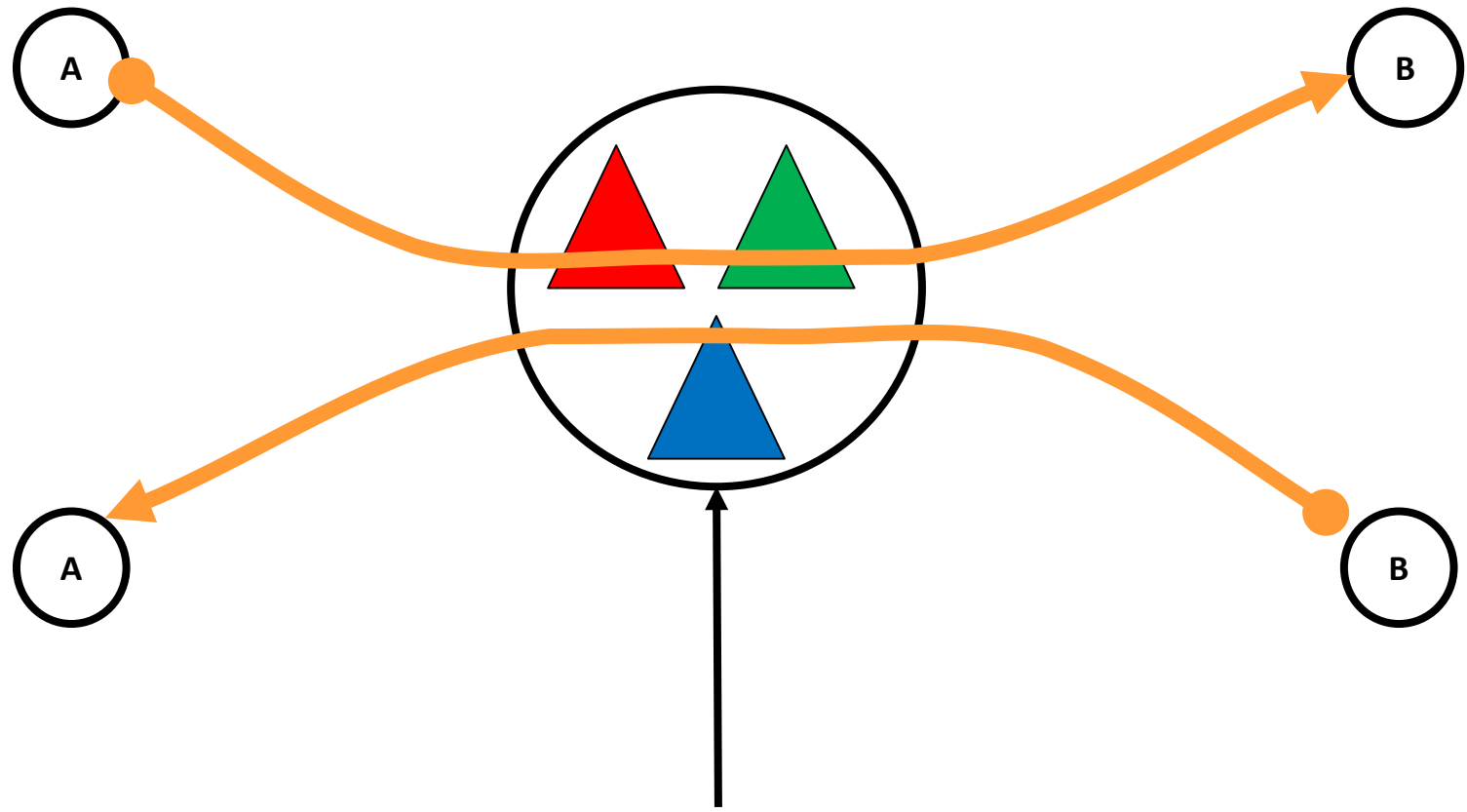
- 1 SC instance per SC leads to suboptimal results
- Having SC instances replicated on every node will lead to optimal results
 - Large capital expenditure to make all nodes NFV capable
 - High Orchestration Overhead for large number of instances
- The question therefore becomes:
 - **How many SC instances to deploy to reduce bandwidth consumption while also reducing nodes used?**
- **We develop a heuristic to help us chose the right number of instances (SPTG)**

Issue of symmetric flows



VNFs placed in different nodes

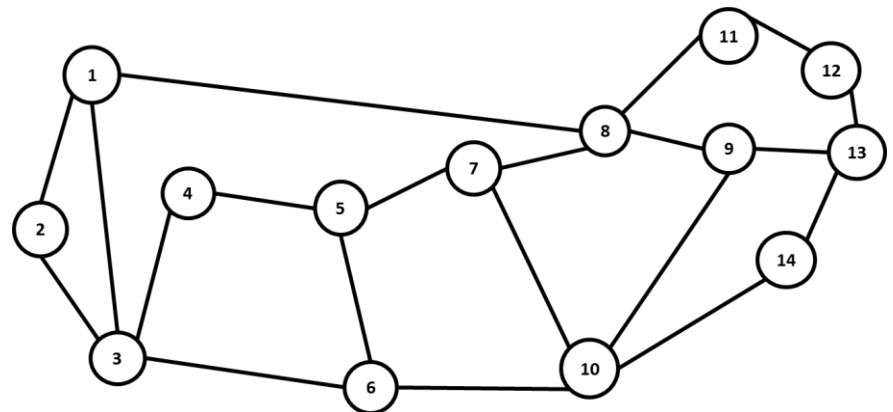
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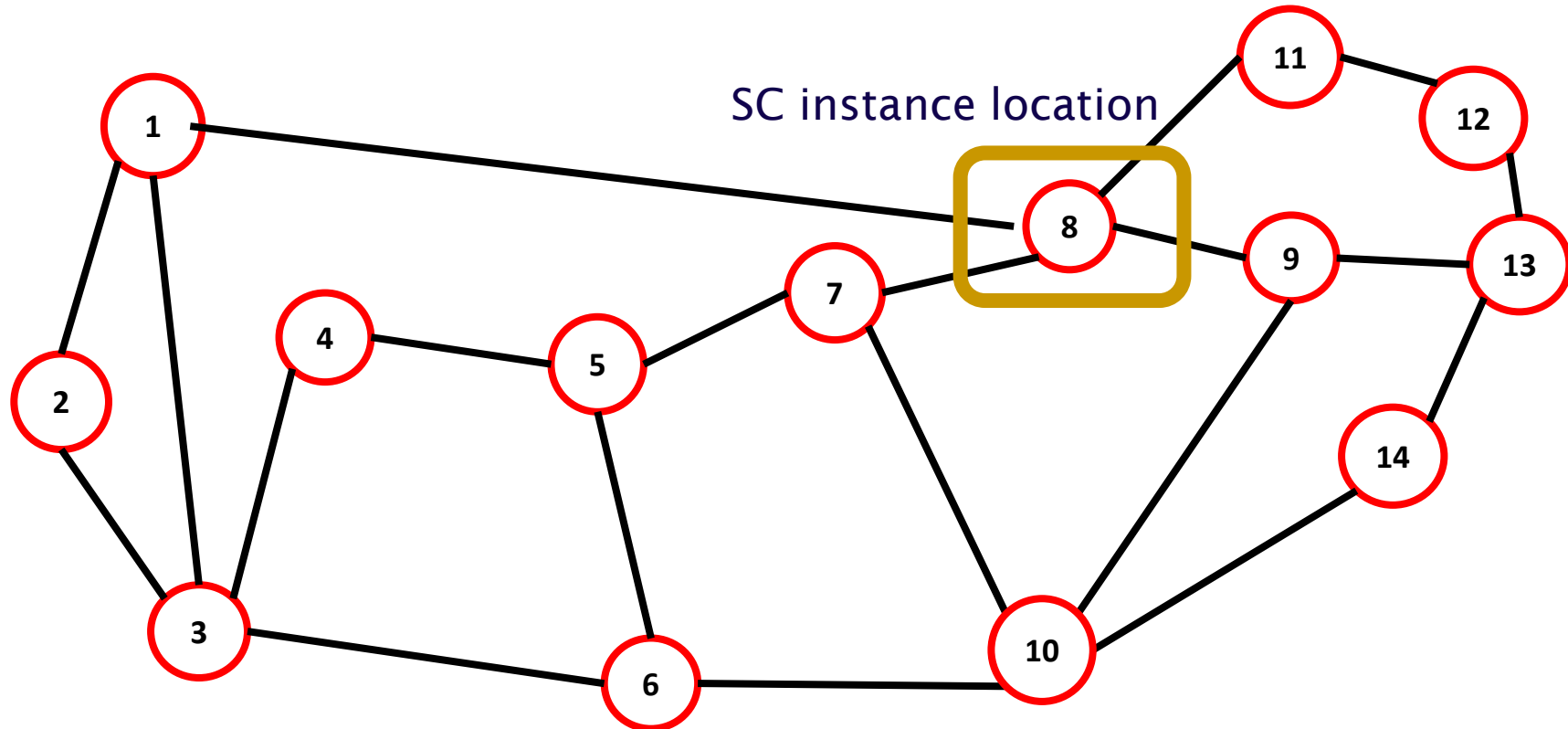
VNFs placed in same node

Continued...

- Placing VNFs for SC at different nodes
 - makes symmetric flow take longer path
- Placing VNFs for SC at one node
 - symmetric flow takes shorter path
 - placement and routing becomes trivial
 - chaining aspect is forgone
 - Is this more realistic?
 - Represents the case of a DC



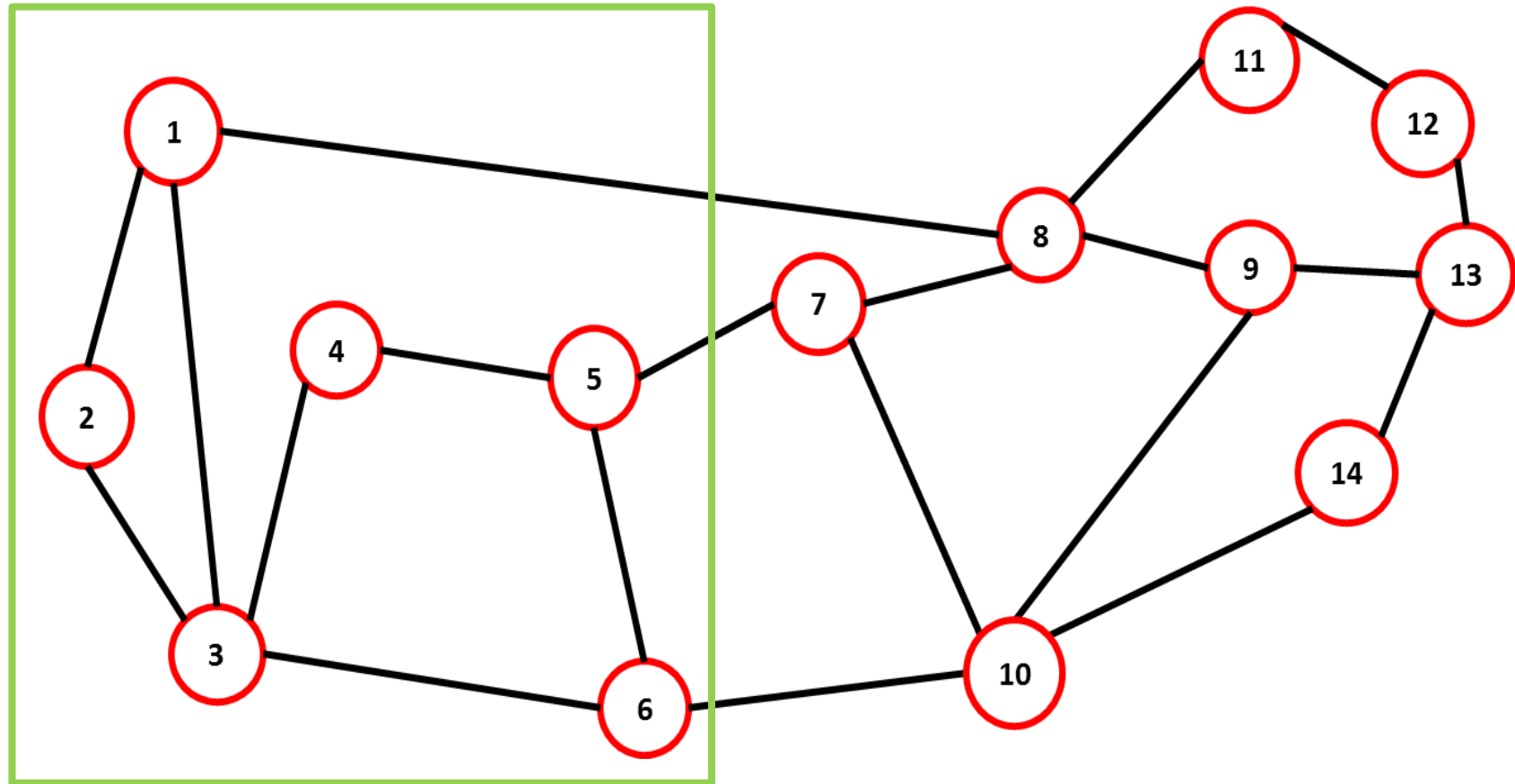
Full Traffic Matrix, 1 SC deployment, 1 SC instance



All nodes are NFV-capable.

All node pairs have requests for the same service chain.

Grouping of traffic pairs



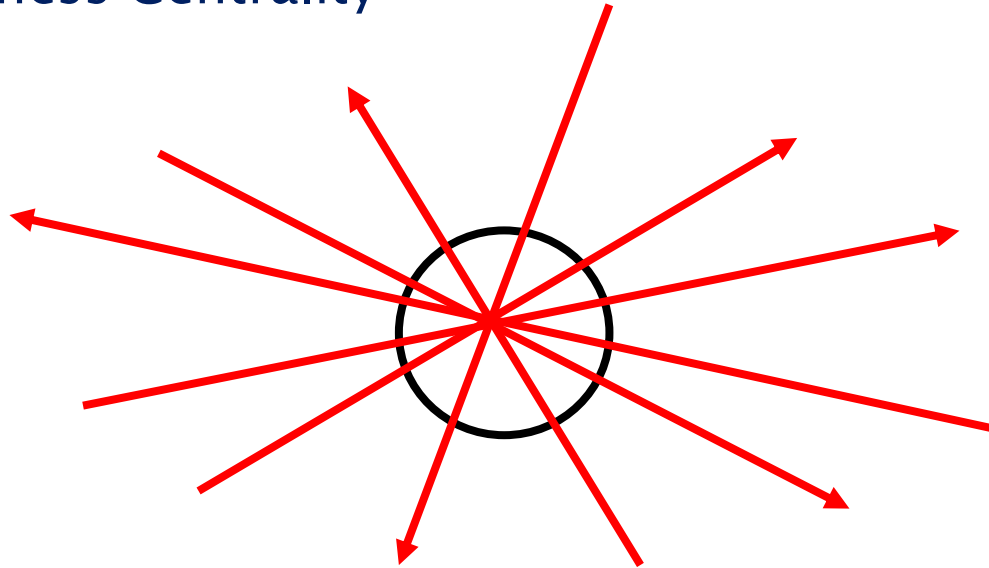
Continued ...



- Create traffic flow groups
- Assign dummy SC Id's to traffic flow groups
- Big Question: **How to do we make traffic groups?**
- Model accounting for traffic groups becomes quadratic. Subsequent, linearization reduced the scalability of the model
- We, therefore, use a heuristic to do make the traffic groups

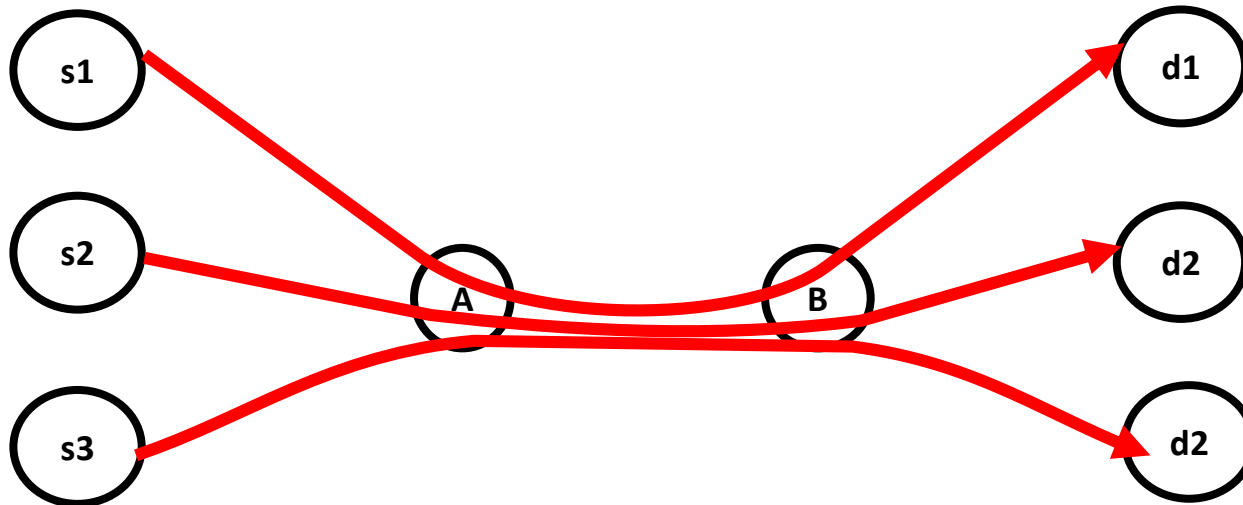
Grouping traffic flows around a node

- Betweenness Centrality



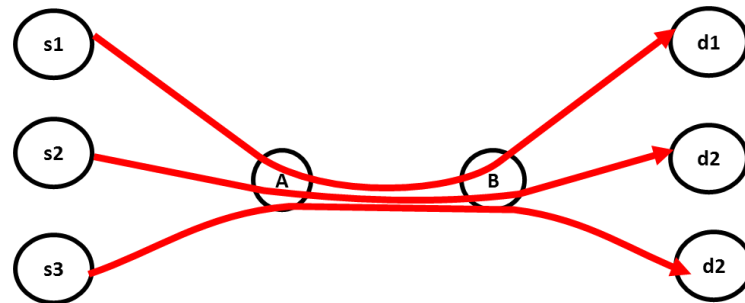
Group around node pairs of the graph

- A and B can also be source and destination
- Done for each SC



Continued...

- Ordered node pair with highest traffic flow count passing through on shortest paths



- Traffic flows which share sub-paths in common
- Deploying one SC instance for each such group

Shortest Path Traffic Grouping (SPTG)

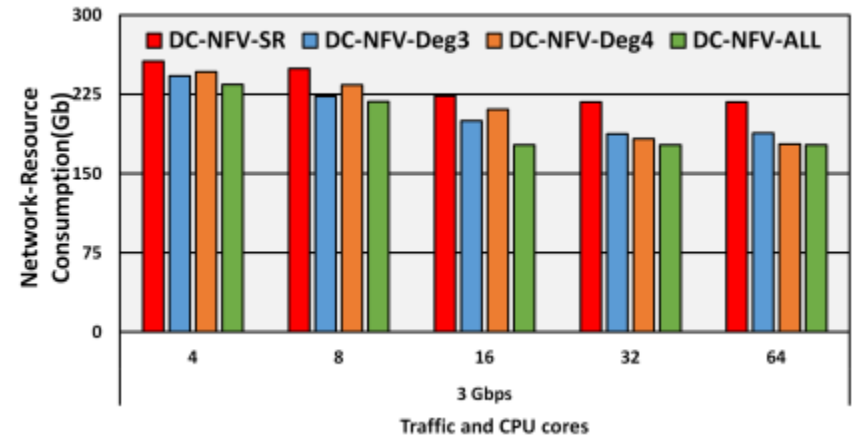
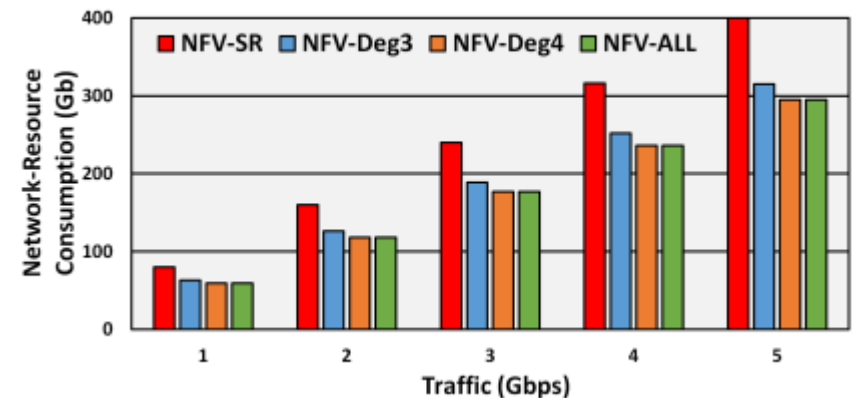
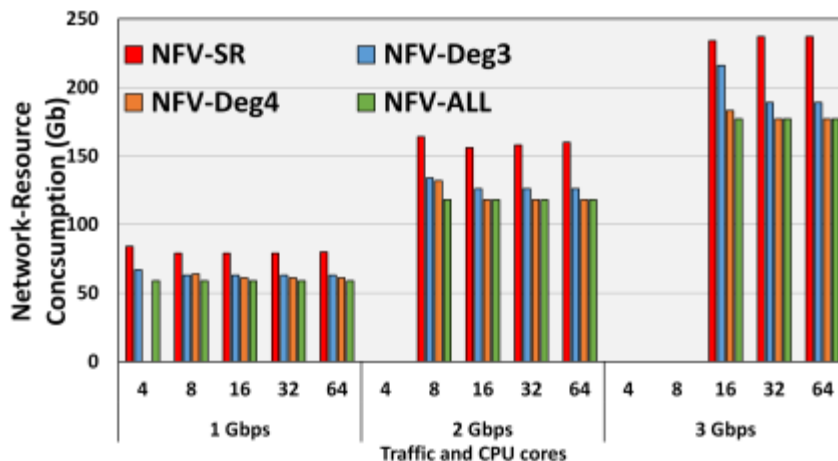
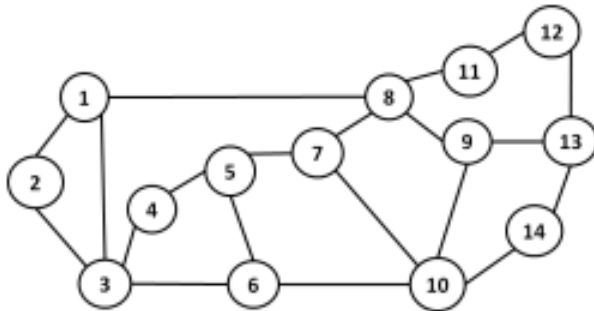
- Given: the number of instances for a SC, all node pairs in a graph G
- The heuristic will:
 1. Find the node pair with the largest number of (s, d) pairs
 2. This becomes another (s, d) pair group
 3. All the (s, d) pairs in the group are removed from the global (s, d) pair list
 4. Repeat step 1 to 3 until number of instances is reached
 5. Iterate through the remaining (s, d) pairs:
 1. Find best group based on which path length through node pair
 2. Add (s, d) pair to that group

2 phase model

- 1st phase
 - Apply SPTG for each SC and create the required number of groups
 - Assign dummy SC ids to groups of (s,d) pairs
- 2nd phase
 - Use the column generation model which decides on 1 SC instance per SC
 - *Also we can control the number of nodes that can host VNFs, we refer to this number as ‘K’*

Previous Results

- 20 uniformly distributed traffic flows, 13 service chains, and 33 VNFs.

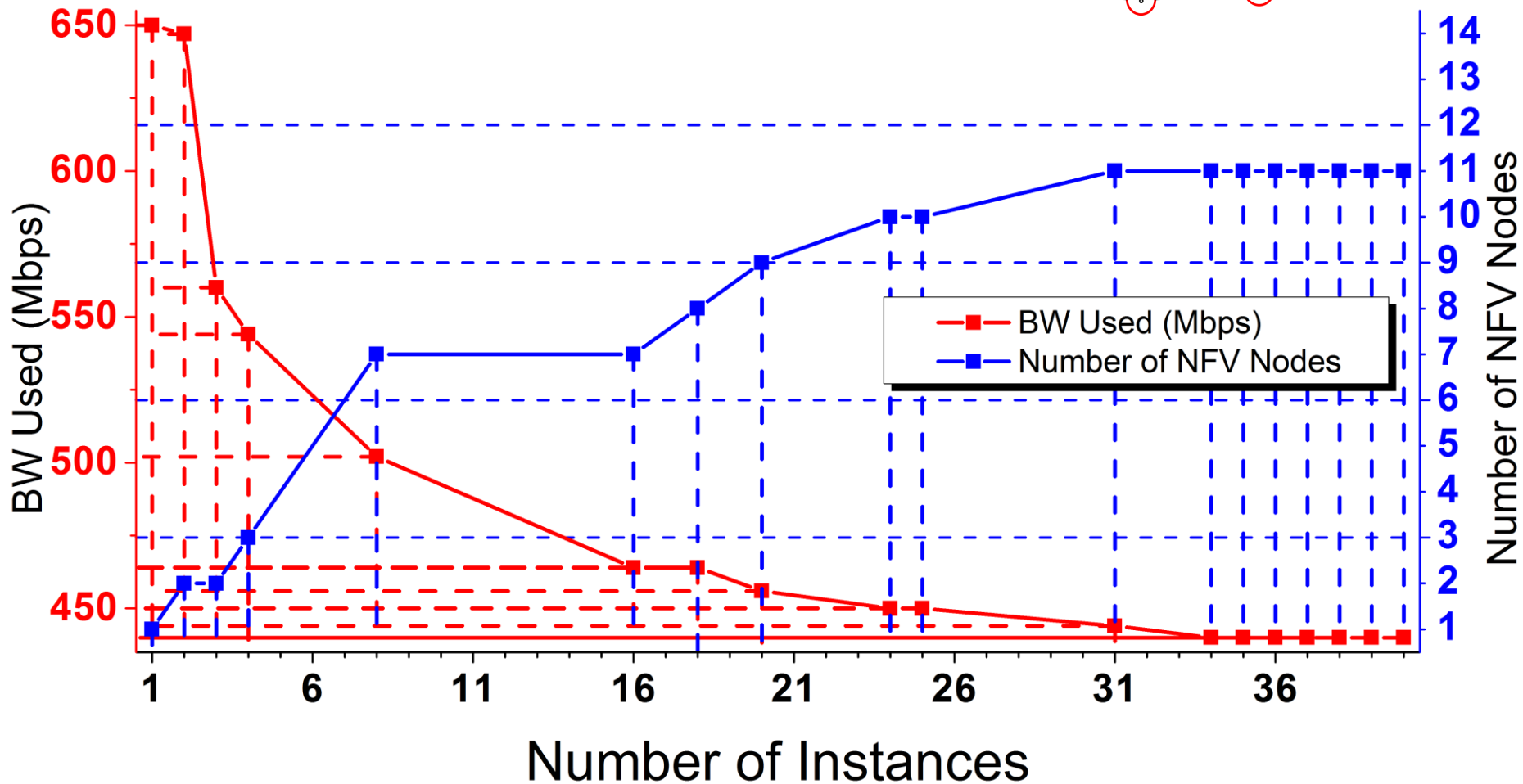
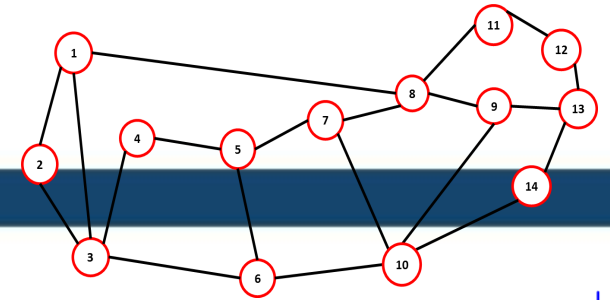


Assumptions

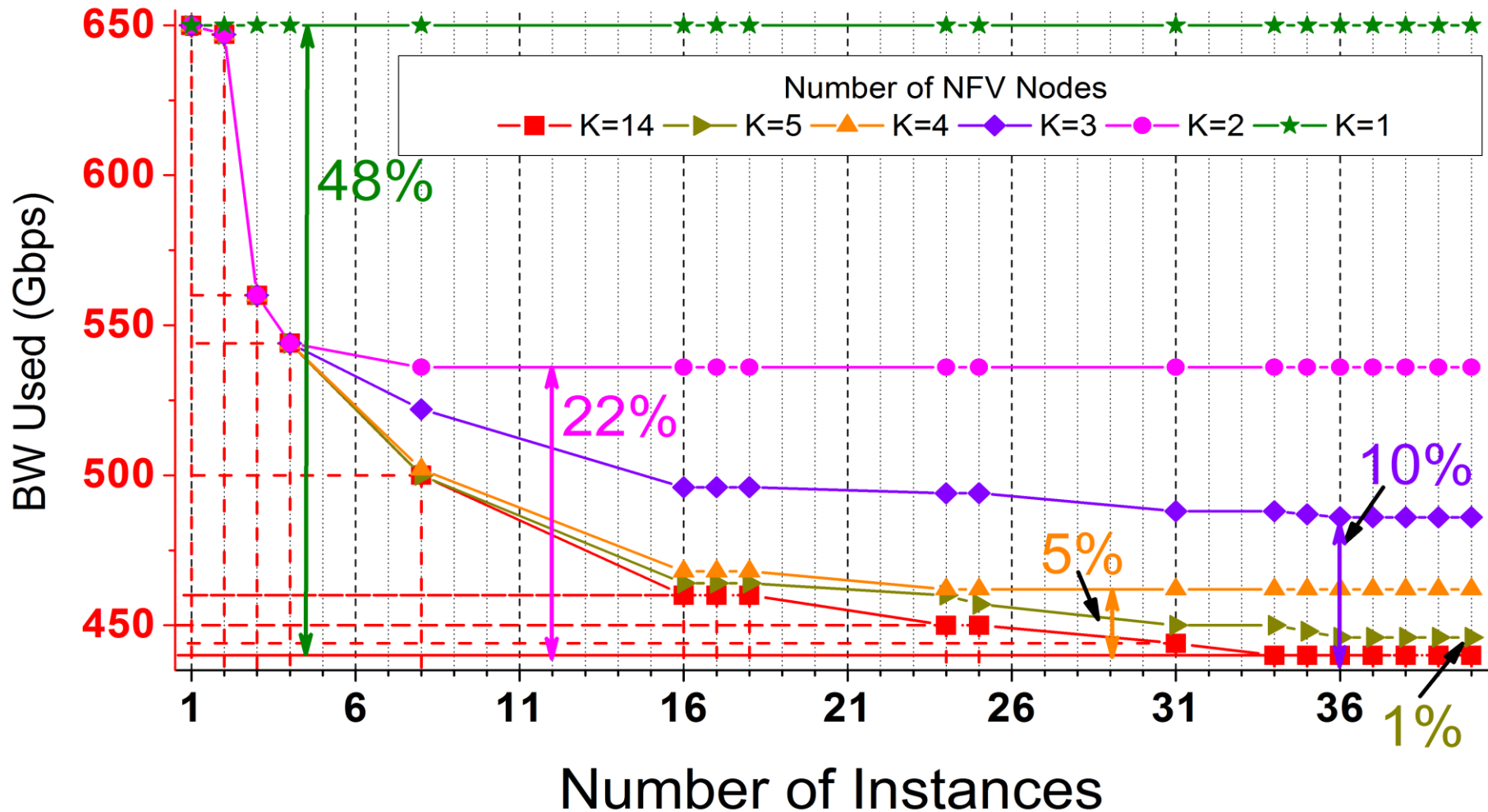
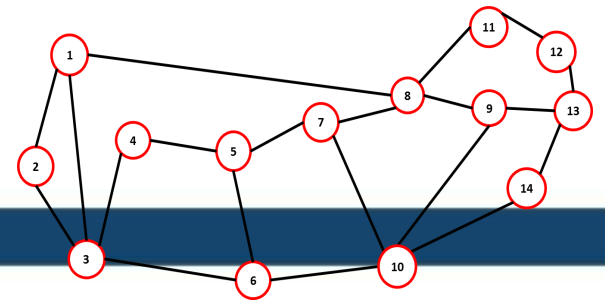


- All nodes are capable of hosting VNFs
- No CPU constraints are enforced
- No link capacity constraints are enforced
- Only one SC instance per SC model
- All traffic pairs have 1 Gb traffic flow

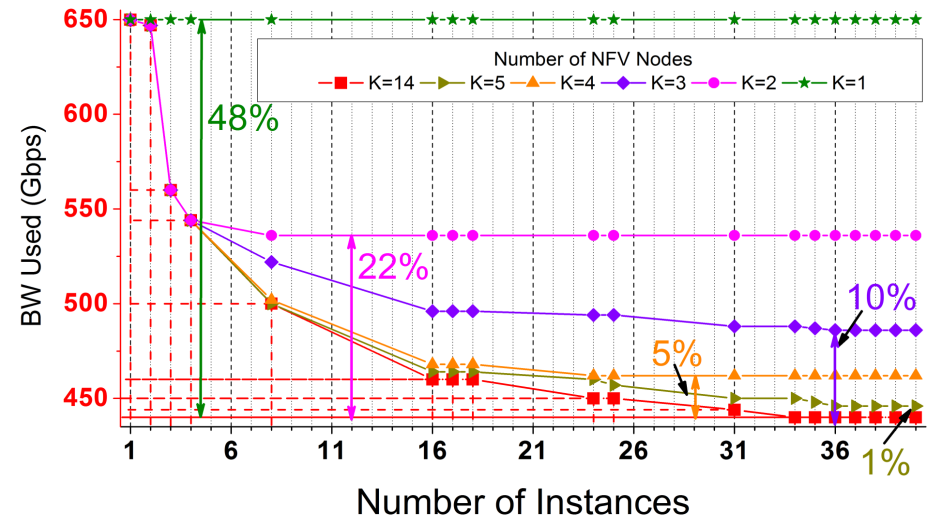
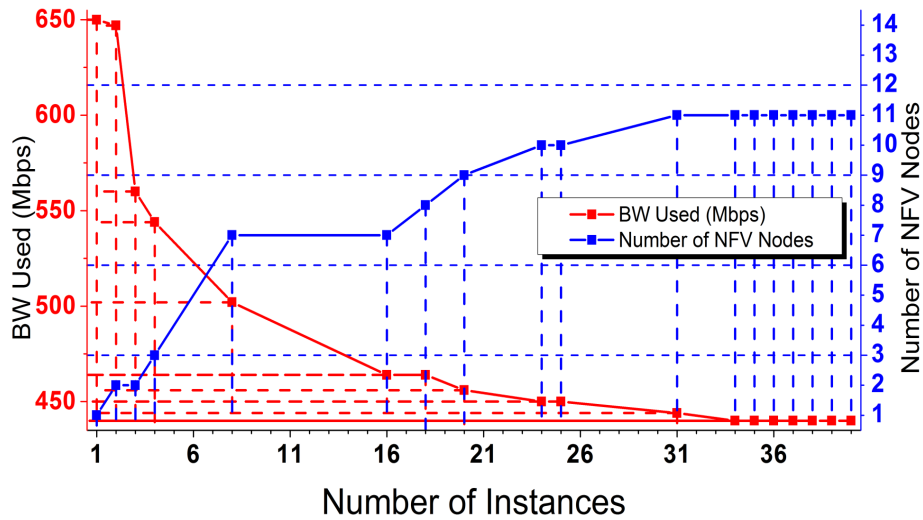
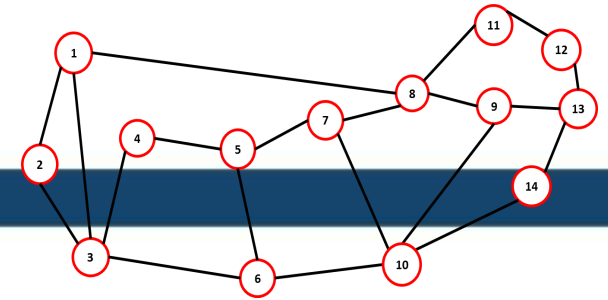
NSFNET K=14 (subset)



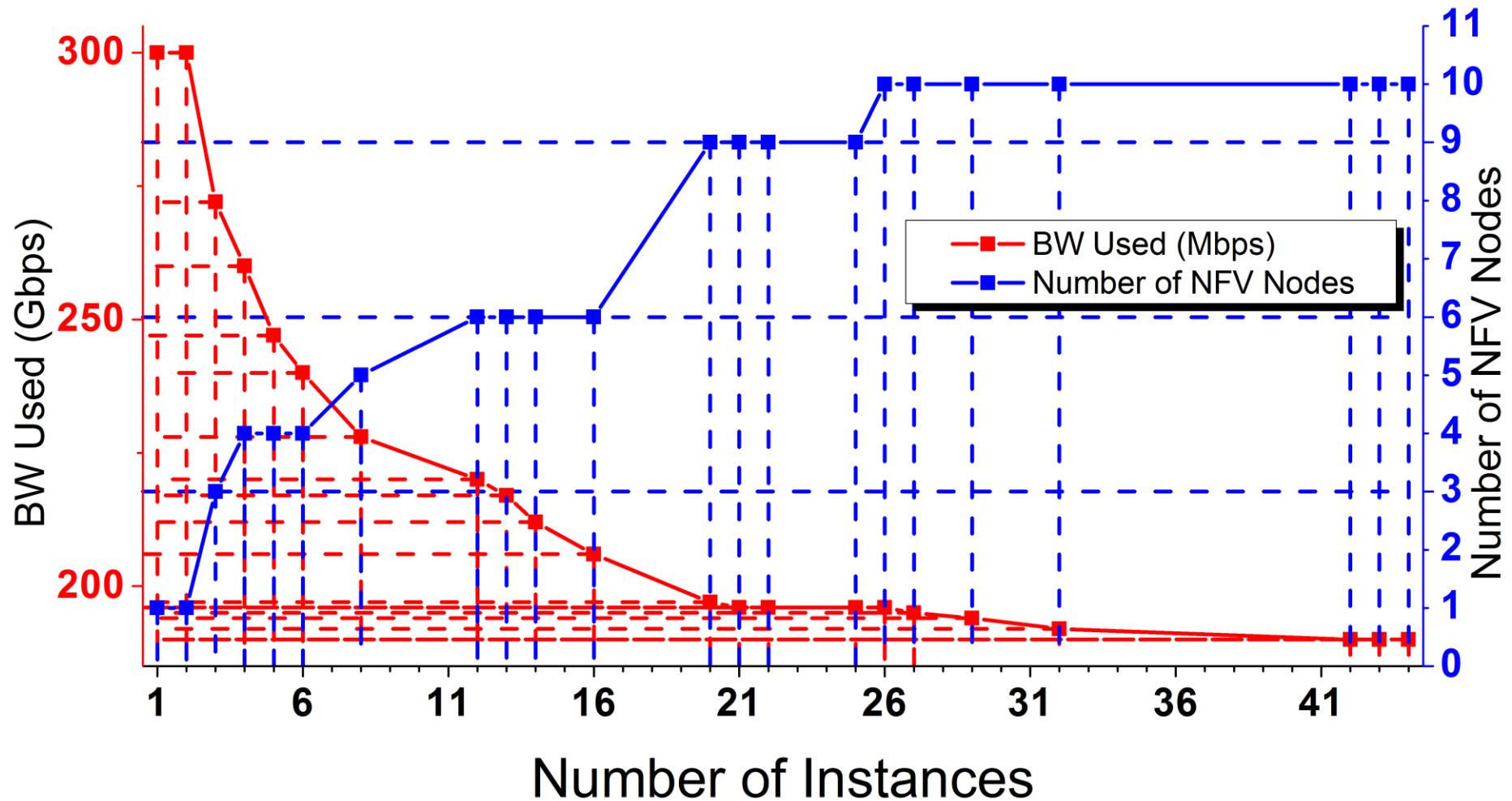
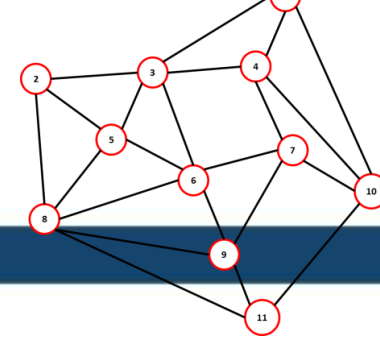
NSFNET K=14,5,4,3,2,1 (subset)



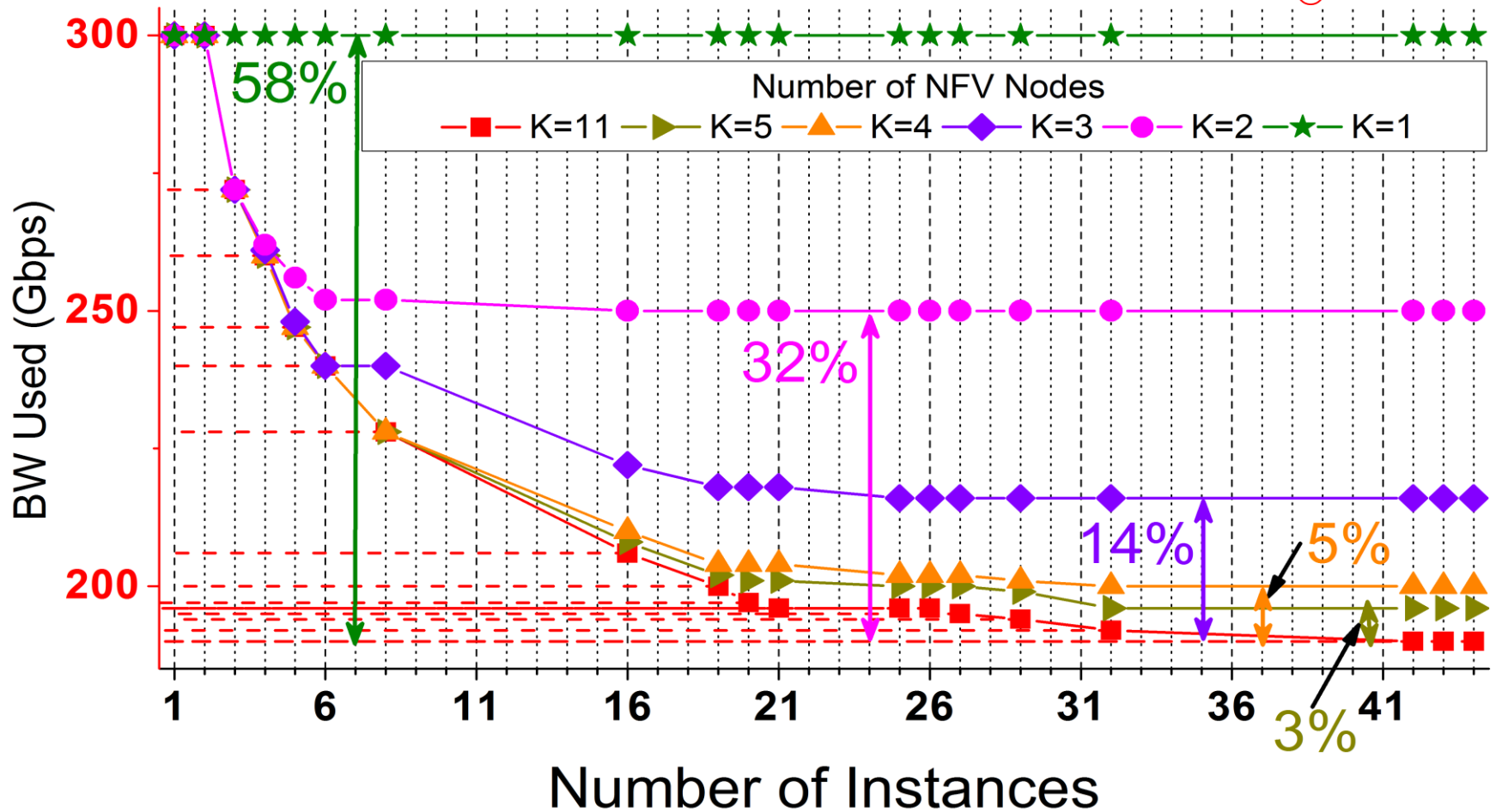
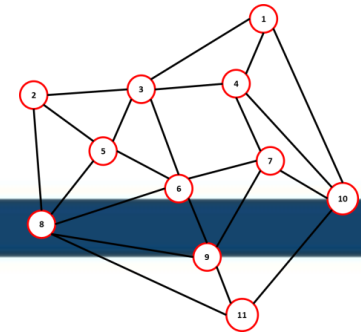
NSFNET



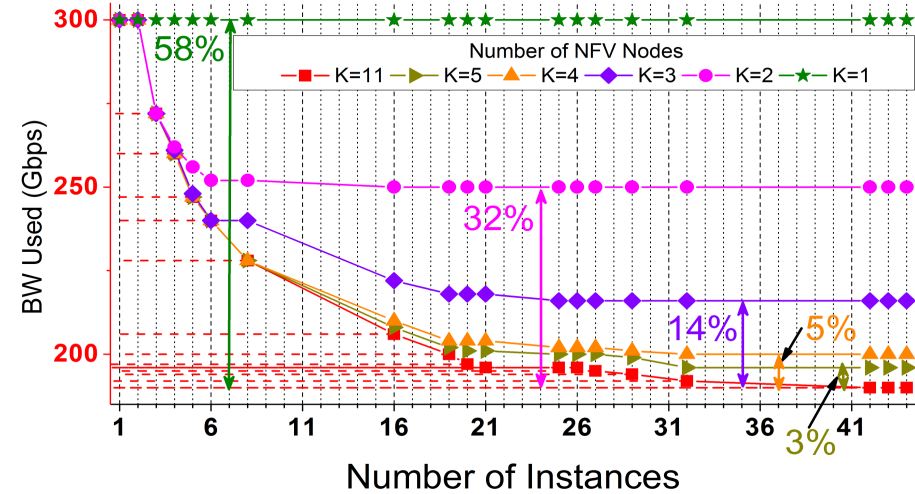
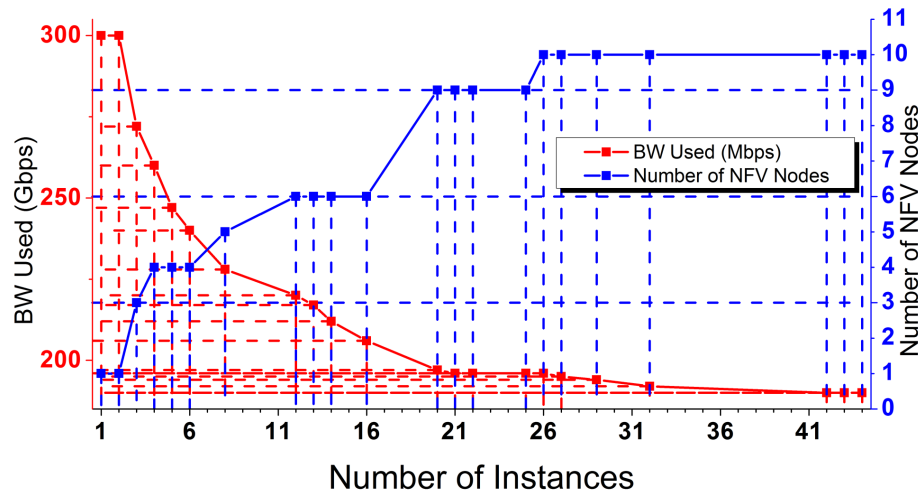
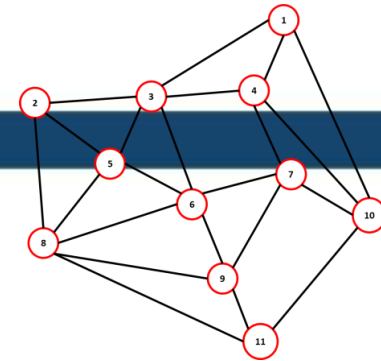
COST239 K=11 (subset)



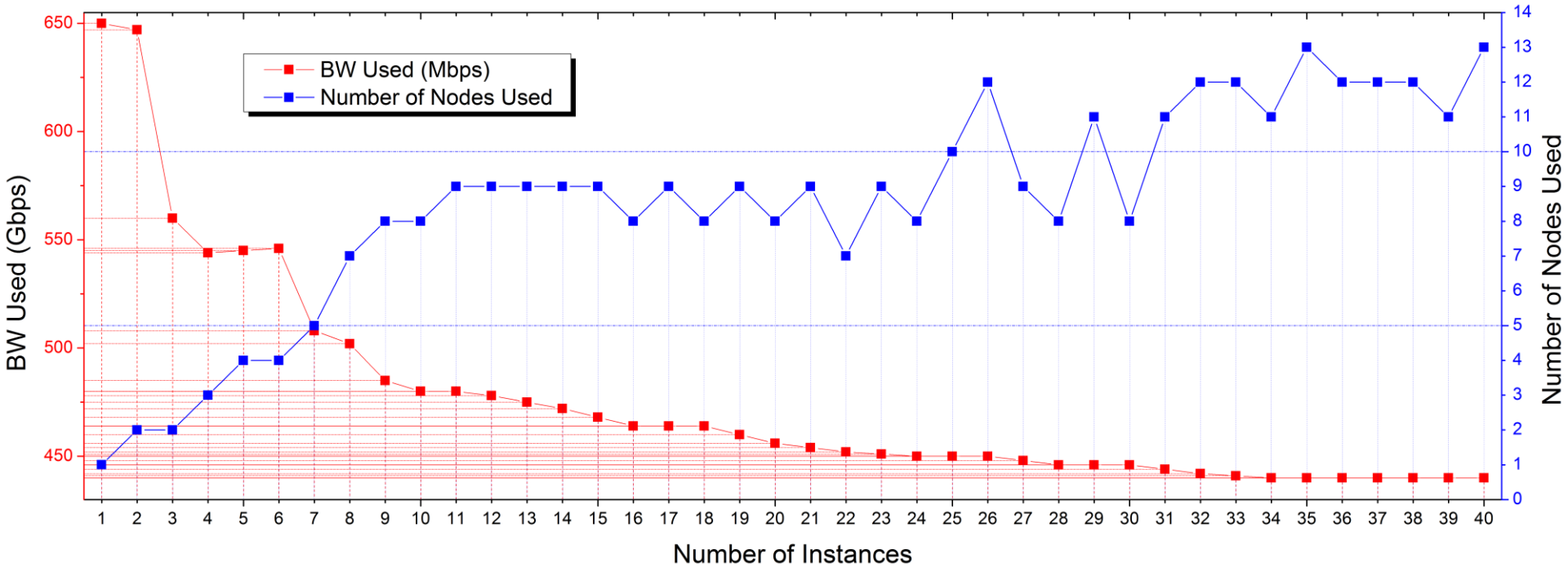
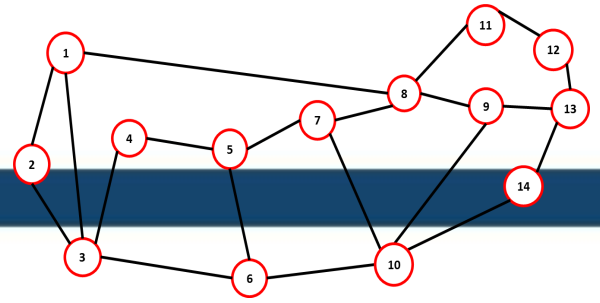
COST239 K=11,5,4,3,2,1 (subset)



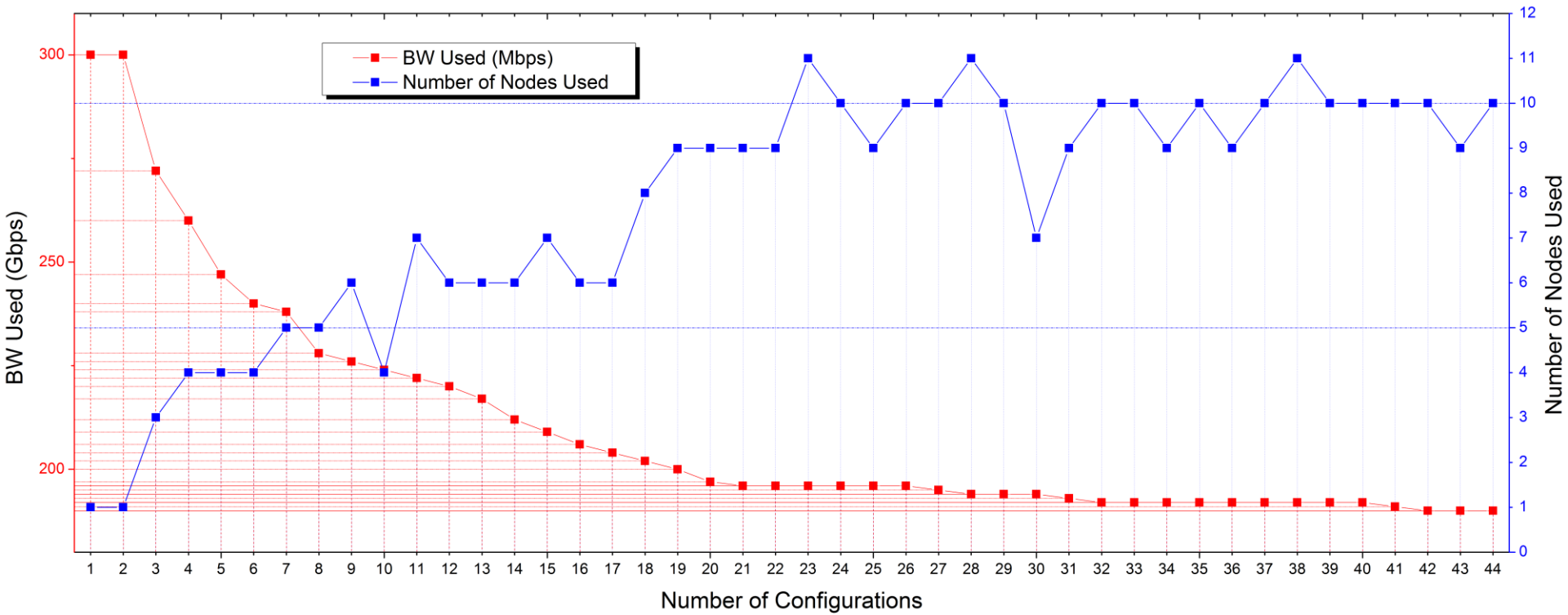
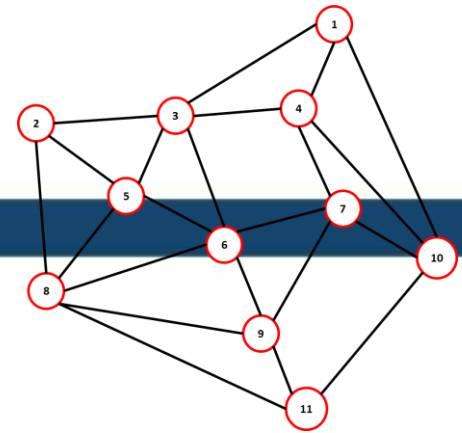
COST239



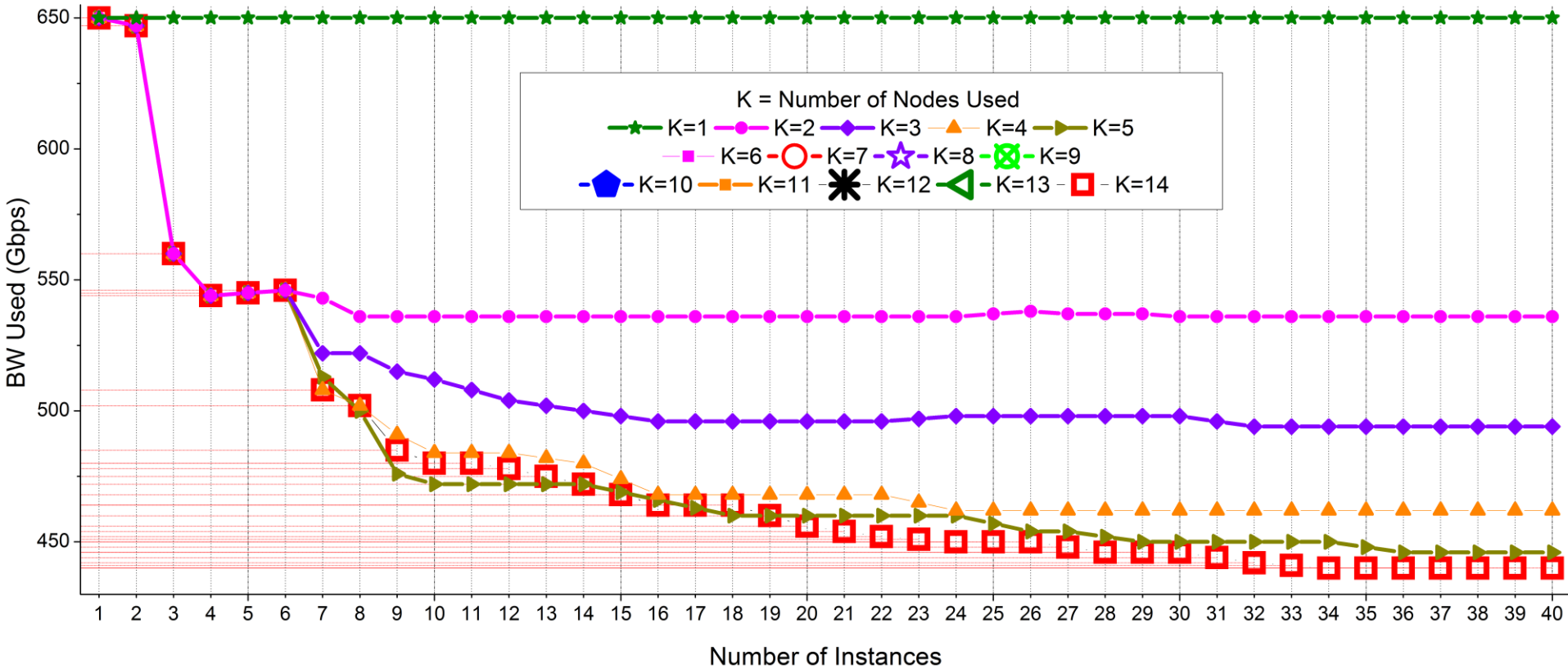
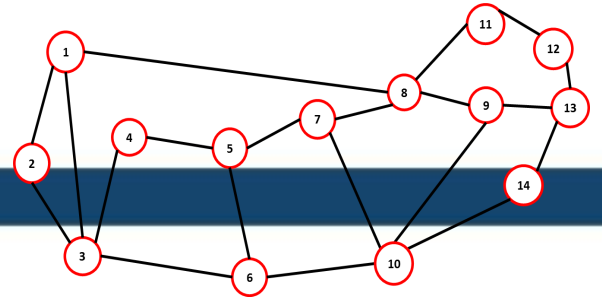
NSFNET K=14



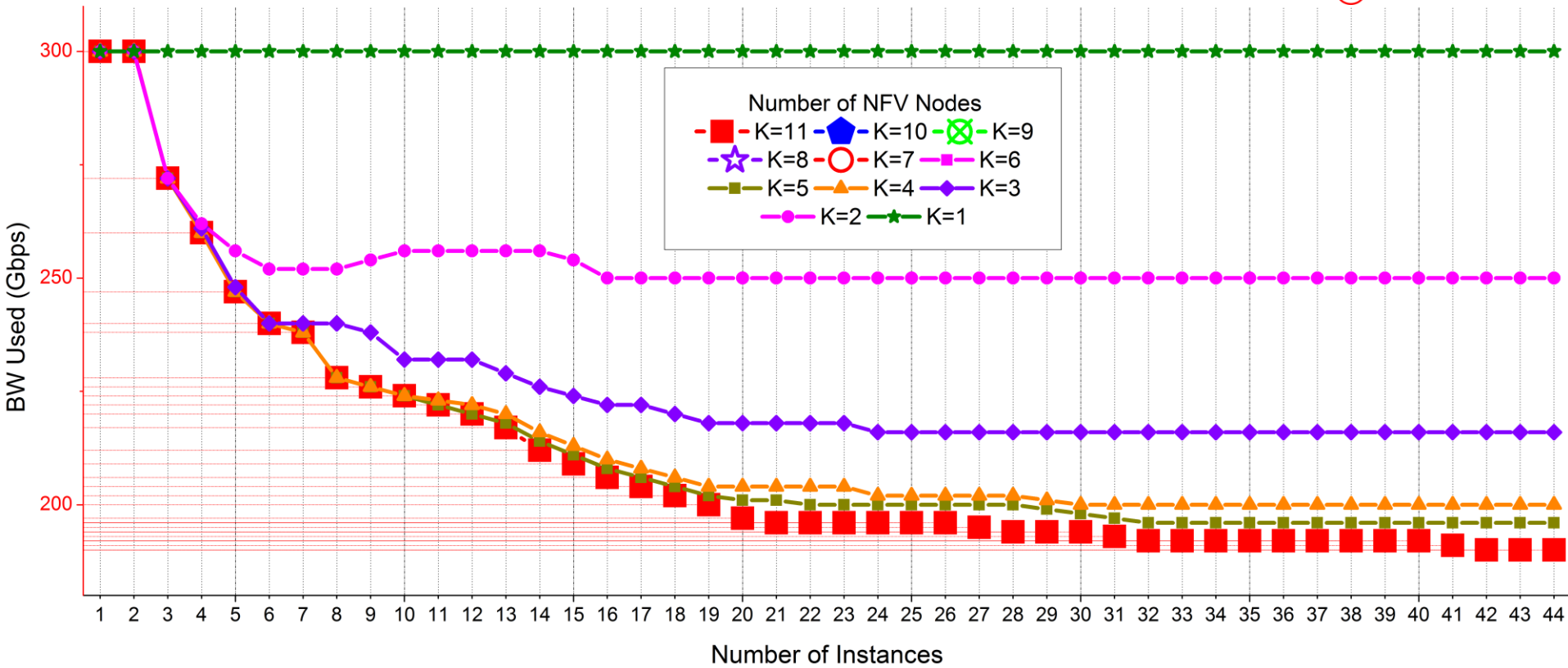
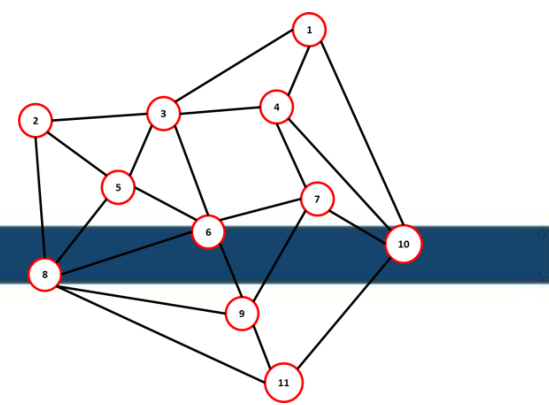
COST239 K=11



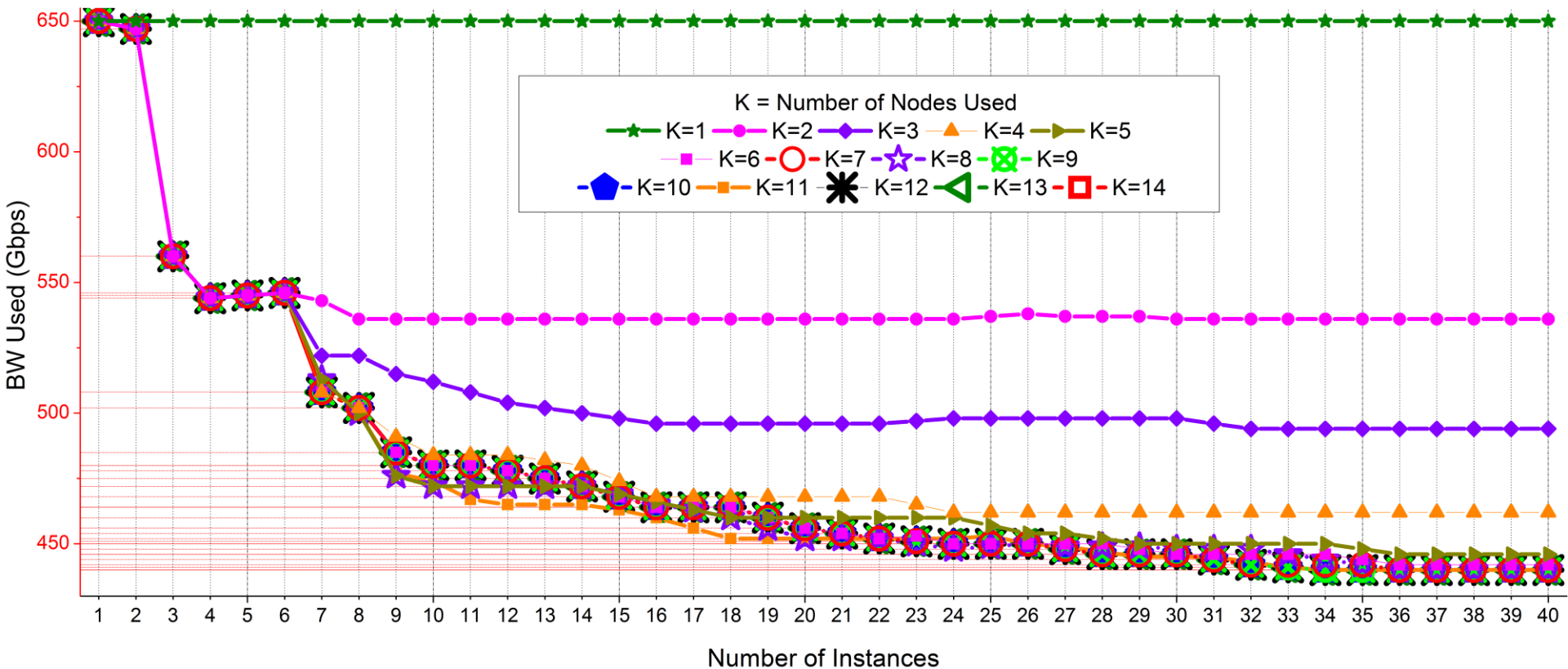
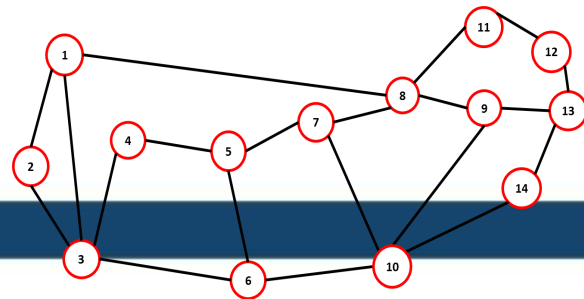
NSFNET K=14,5,4,3,2,1



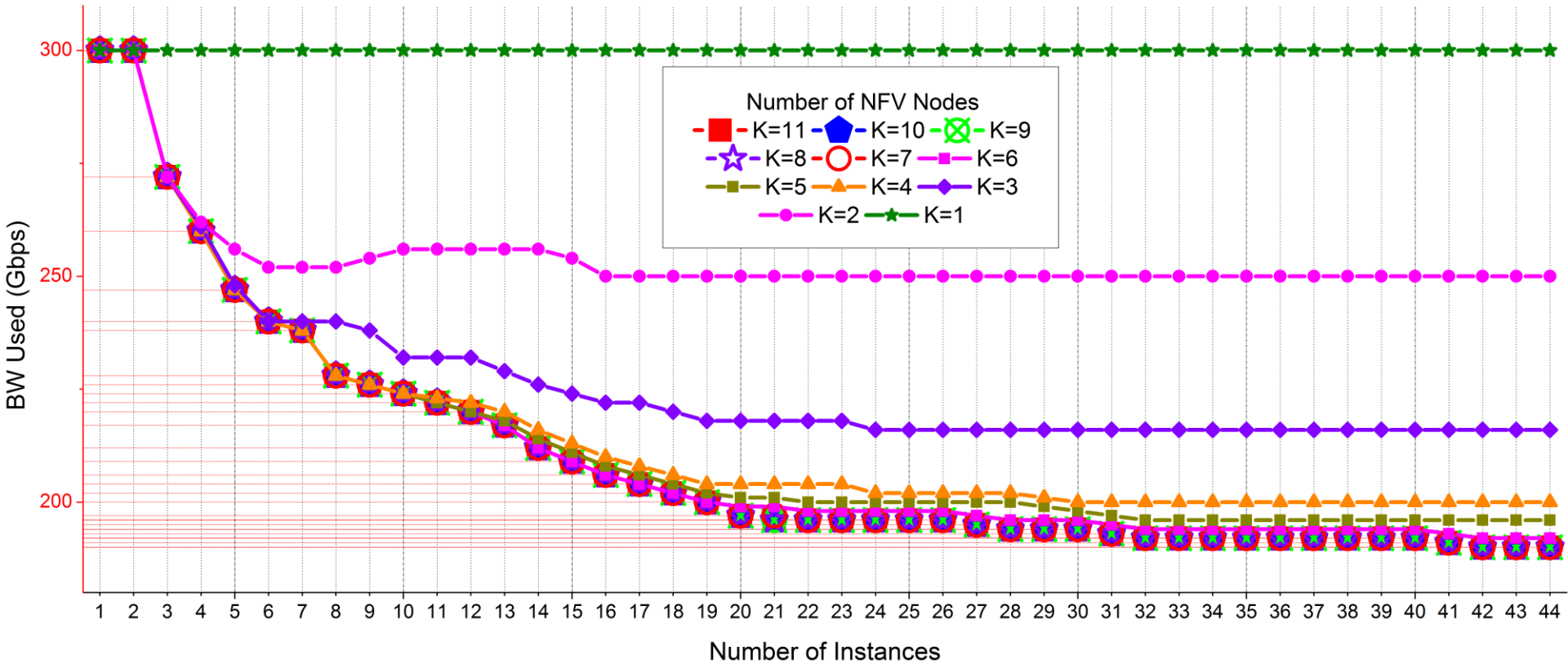
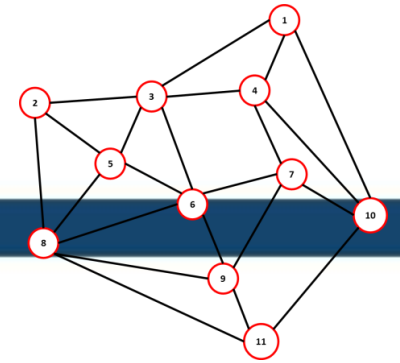
COST239 K=11,5,4,3,2,1



NSFNET K=ALL



COST239 K=ALL



Future Work Directions



- **Comparison of 2 phase approach with quadratic model**
- **Set of results for when CPU and link capacity constraints are enforced:**
 - Difference from current results
- **Cases where distribution of VNFs occur:**
 - Cases where CPU resources are constrained or VNF replicas (because of licenses) are enforced
 - Any additional cases?

Continued...

- **Scalability of the approach for larger network instances**
- **Account for non-uniform traffic :**
 - Matrices which are still dense but with a mixture of large and small flows. How to form traffic flow groups then?
 - Also account for multiple heterogeneous service chains
- Can the current group of traffic flows be improved?
- General graph VNFs instead of linear service chain?