## Advanced Disaster-Resiliency Strategies for Next-Generation Metro Optical Networks in the Context of Smart Cities

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## Outline

- 1. Chapter 1: Background
- 2. Chapter 2: Logical network mapping with content connectivity against multiple link failures in metro optical networks
- 3. Chapter 3: Ongoing and future research



# Chapter 1: Background



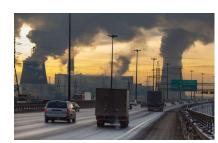
#### **Urbanization**

**Pollutions:** 

#### Population trend:

## World urban population: 55% in 2018, 68% in 2050 (expected)

• In U.S., urban population: **82**% in 2018 (Source: United Nations)



Source: medicaldaily.com

#### **Energy crisis:**



Source: Indianfolk.com

#### **Issues**:

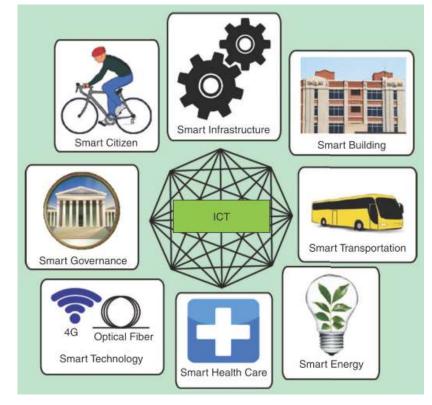
- Uncontrolled growth of urban population
- Limited natural and man-made resources

**Solution**: Smart cities



## Smart Cities: Definition and Components

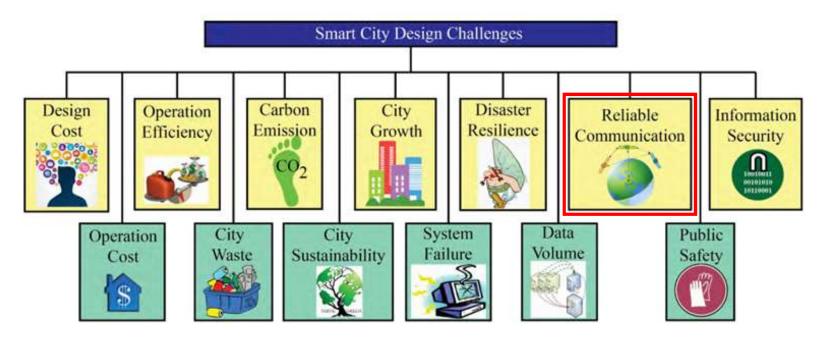
- Smart cities: effective approach to manage limited resources to serve largest possible population in order to improve:
  - ✓ Livability,
  - √ Workability,
  - ✓ and Sustainability [1]



Smart city components [1]



## Smart Cities: Design Challenges



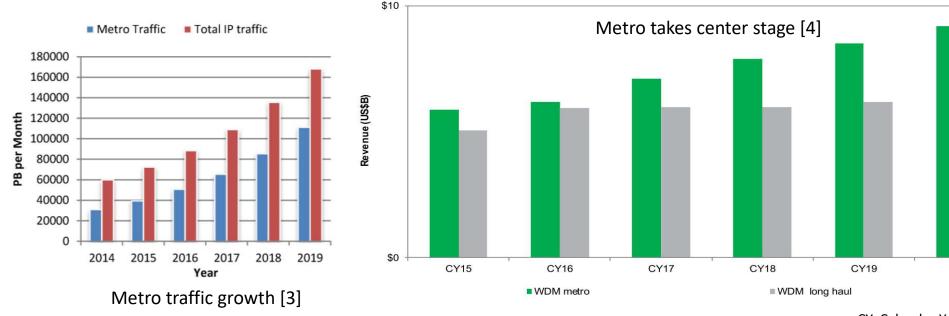
#### Reliable communication as one of design challenges [1], [2]



<sup>[1]</sup> S. P. Mohanty, U. Choppali, and E. Kougianos, "Everything you wanted to know about smart cities: The Internet of things is the backbone," *IEEE Consumer Electronics Magazine*, vol. 5, no. 3, pp. 60-70, July 2016.

<sup>[2]</sup> S. A. Shah, D. Z. Seker, M. M. Rathore, S. Hameed, S. Ben Yahia, and D. Draheim, "Towards Disaster Resilient Smart Cities: Can Internet of Things and Big Data Analytics Be the Game Changers?," *IEEE Access*, vol. 7, pp. 91885-91903, 2019.

## Increasing Role of Metro Optical Networks



CY: Calendar Year

- 66% total IP traffic in 2019 [3]
- Metro hardware revenue surpasses long haul's [4]



CY20

## **Emerging Services**





#### More stringent requirements for:



- ✓ Latency
- ✓ Reliability: For example, uRLLC in 5G requires 99.999% (five nines) availability [5]





Reliability in metro optical networks is gaining importance, particularly in the context of smart cities

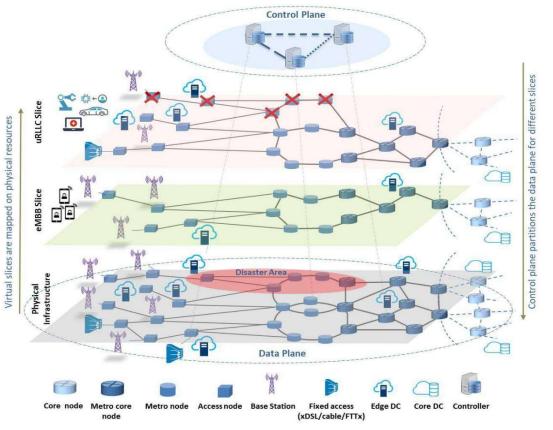




**New services** 



## Reliability in Next-Generation Metro Optical Networks



#### **Technology enablers:**

- ✓ Software-Defined Networking
- ✓ Network Functions Virtualization
- ✓ Network Slicing

#### **Disasters may disrupt services**

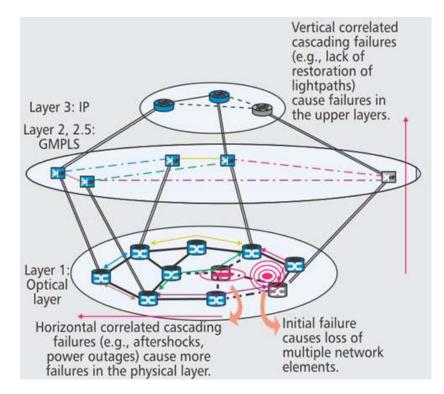
Resiliency strategies must be adapted to the new architecture and enablers

Next-Generation Metro Optical Network Architecture [6]



#### Disaster Failure Characteristics

- Disruption of multiple links and nodes
- Cascading:
  - √ horizontal (optical layer)
  - ✓ vertical (higher layers)
- We focus on link failures (higher probability [8])



#### Cascading failures in optical networks [7]

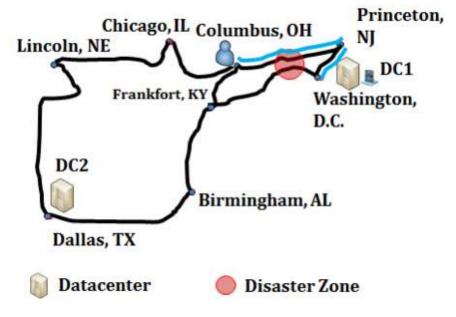


<sup>[7]</sup> B. Mukherjee, M. F. Habib, and F. Dikbiyik, "Network adaptability from disaster disruptions and cascading failures," *IEEE Communications Magazine*, vol. 52, no. 5, pp. 230-238, May 2014.

<sup>[8]</sup> P. Gill, N. Jain, and N. Nagappan, "Understanding network failures in data centers: measurement, analysis, and implications," *Proceedings of ACM SIGCOMM*, vol. 41, no. 4, Aug. 2011.

## Network Connectivity (NC)

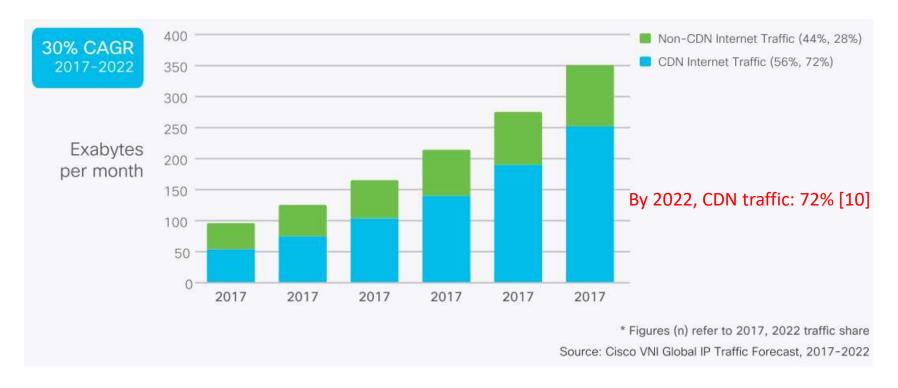
- Reachability of every network node from all other nodes
- Default metric for network survivability
- In case of disasters, NC may not be possible



Scenario where NC is not possible [9]



## Increasing Role of Content Delivery Networks (CDN)



#### **Content connectivity -> Service continuity (most applications)**



## **Content Connectivity**

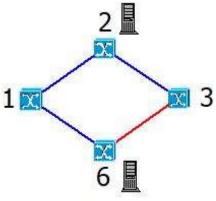
- ✓ Reachability of content from every node in a network under a given failure scenario [9]
- ✓ Important survivability metric
- ✓ Possible in some scenarios NC impossible



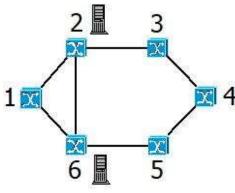
Content Connectivity is guaranteed after disaster [9]



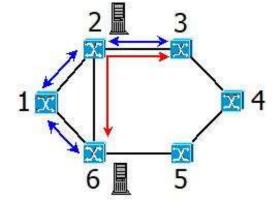
## Mapping With Content Connectivity



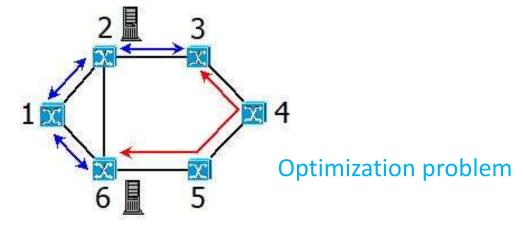
Logical topology



Physical topology



2-3 link failure disconnects 3 from content



Content-connected against single link failures



## Chapter 2:

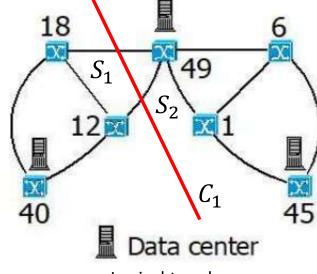
# Logical Network Mapping With Content Connectivity Against Multiple Link Failures in Metro Optical Networks

(Preliminary results in this chapter have been submitted to ANTS 2019)



## Cut and Cutset of a Network

- Cut: partition of the network into two disconnected segments (e.g., two node groups  $S_1$  and  $S_2$ )
- Cutset: set of links with one endpoint in  $S_1$  and the other in  $S_2$



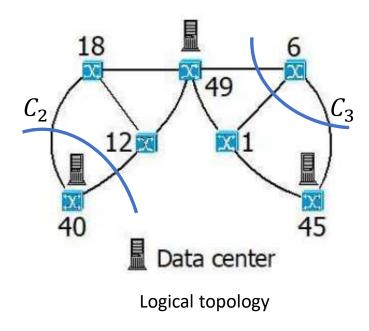
Logical topology

- $\checkmark$   $C_1$  is a cut
- $\checkmark S_1 = \{12, 18, 40\}$
- $\checkmark S_2 = \{1, 6, 45, 49\}$
- ✓ {18-49, 12-49} is a cutset



#### Network Cutset and Content Cutset

- Network Connectivity (NC) cutset:
  - $\checkmark C_2$
  - ✓ Removal all links in  $C_2$  violates NC
- Content Connectivity (CC) cutset:
  - $\checkmark C_3$
  - ✓ Removal all links in  $C_3$  disconnects node 6 from content
  - ✓ Nodes co-located with datacenters are content-connected



Content available at all DCs



## **Problem Statement**

#### Given:

- ✓ Logical topology
- √ Physical topology

#### Objective:

✓ Minimize network resource usage

#### Output:

✓ Mapping with content connectivity after n link failures



#### **Problem Notations**

- NC -n: Network Connectivity (NC) after failures on n physical links
- CC -n: Content Connectivity (NC) after failures on n physical links



## Input Parameters

- $G_P(V_P, E_P)$ : physical topology (graph)
- $V_P$ : set of physical nodes
- $E_P$ : set of physical links
- $G_L(V_L, E_L)$ : logical topology (graph)
- $V_L$ : set of logical nodes
- $E_L$ : set of logical links

- D: set of Datacenter,  $D \subset V_L$
- $F_{ij}$ : number of fiber from i to j
- *W*: number of wavelength/fiber
- *n*: number of physical link failures
- $P_n$ : set of n physical links
- C<sub>cc</sub>: set of content-connected cutsets (next slides)



## $P_n$ : Set of n Physical Links

- Number of physical nodes:  $N_P = |V_P|$ , where |. | set cardinality
- Number of physical links:  $L_P = |E_P|$
- Select n links out of  $L_P$  links: Combination without order and repetition
- Total number of valid sets:

$$\frac{L_P!}{n! (L_P - n)!}$$



## $C_{CC}$ : Set of Content-Connected Cutsets

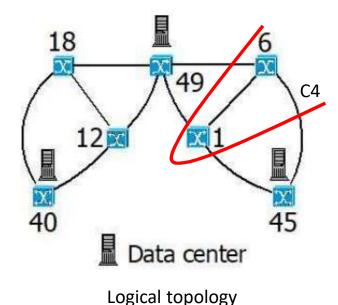
- Number of logical nodes:  $N_L = |V_L|$
- Number of logical links:  $L_L = |E_L|$
- Content-Connected (CC) cutset:
  - ✓ C4 (for example)

$$\checkmark S_1 = \{1, 6\}$$

$$\checkmark S_2 = V_L - S_1 = \{12, 18, 40, 45, 49\}$$

$$\checkmark S_1 \cap D = \{40, 45, 49\} = \emptyset$$

•  $C_{CC}$ : enumeration of all CC cutsets



For example, this logical topology has 30 CC cutsets



#### CC-*n* Existence

Theorem 1: Given  $G_P(V_P, E_P)$ ,  $G_L(V_L, E_L)$ , and D, to find the mapping of  $G_L$  over  $G_P$  that guarantees CC-n, the following conditions must be satisfied:

- ✓ each logical node  $s \in V_L D$  has a nodal degree  $\delta(s) \ge n + 1$ , and
- $\checkmark$  each physical node  $i \in V_P$ : i = s has a nodal degree  $\delta(i) \ge n + 1$ .

logical link S logical link

physical link i physical link

 $\delta(s)=2$ , no CC-2 solution for node s  $\delta(i)=2$ , no CC-2 solution mapped over node i

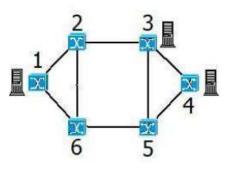


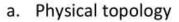
## Problem Variable

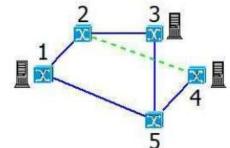
• Variable definition (binary):

 $\checkmark f_{ij}^{st} = 1$  if logical link st is mapped over physical link ij

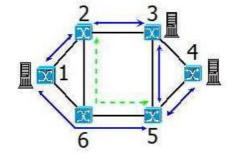
 $\checkmark f_{ij}^{st} = 0$  otherwise







b. Logical topology



c. Mapping

For example:  $f_{16}^{15} = f_{65}^{15} = 1$ ,  $f_{ij}^{15} = 0$ ,  $\forall ij \notin \{16,65\}$ 



#### CC-*n* Enforcement

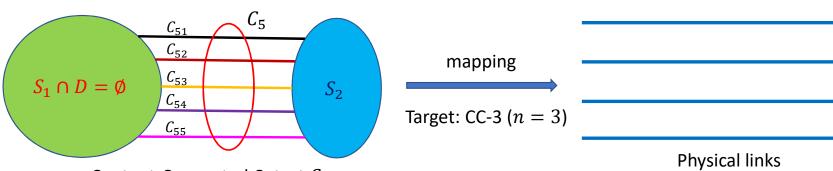
Theorem 2: Given  $G_P(V_P, E_P)$ ,  $G_L(V_L, E_L)$ , D, let  $P_n = \{\{P_n^k\}: |\{P_n^k\}| = n, \{P_n^k\} \in E_P\}$  be set of all possible combinations of n distinct physical links, and  $C_{CC} = \{C_{CC}^l(S_l, V_L - S_l): S_l \cap D = \emptyset\}$  be set of logical topology content-connected cutsets where the removal of all logical links in each cutset  $C_{CC}^l$  disconnects  $G_L$  and divides  $V_L$  into two disjoint sets with one set without datacenters, the mapping of  $G_L$  over  $G_P$  is CC-n if and only if:

$$\sum_{ij\in P_n^k, st\in C_{CC}^l} f_{ij}^{st} \leq \left| C_{CC}^l \right| - 1, \forall P_n^k \in P_n, \forall C_{CC}^l \in C_{CC}.$$

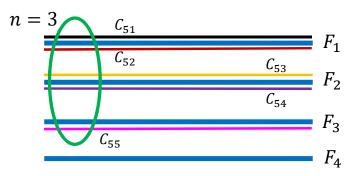


### CC-*n* Enforcement

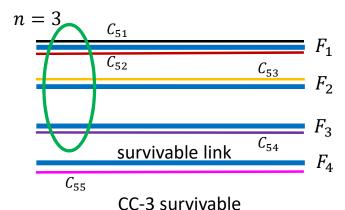
### Theorem 2: Example n=3 (survivable against 3 link failures)



Content-Connected Cutset  $C_5$ 



CC-3 not survivable





 $F_2$ 

 $F_4$ 

## Mathematical Formulations of CC-n Problem

#### *Objective function:*

$$\min \sum_{ij \in E_P, st \in E_L} f_{ij}^{st}$$

- ✓ Result in an ILP
- ✓ Lower complexity

Capacity Constr.

jective function: Subject to: Capacity Const  $\min \sum_{i=1}^{n} f_{ij}^{st} \qquad \sum_{st \in E_L} f_{ij}^{st} \leq F_{ij} \times W, \forall ij \in E_P$ 

Flow Constr.

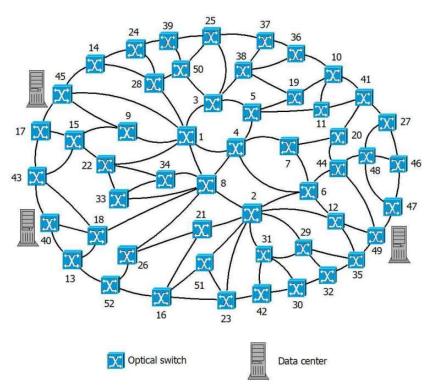
• 
$$\sum_{j:ji\in E_P} f_{ji}^{st} - \sum_{j:ij\in E_P} f_{ij}^{st} = \begin{cases} -1 \text{ if } i = s \\ 1 \text{ if } i = t \\ 0 \text{ otherwise} \end{cases}$$
  
 $\forall i \in V_P, \forall st \in E_L$ 

CC-n Constr.

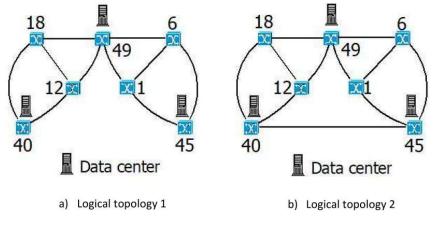
$$\sum_{\substack{ij \in P_n^k, st \in C_{cc}^l \\ \forall P_n^k \in P_n, \forall C_{cc}^l \in C_{cc}}} f_{ij}^{st} \leq \left| C_{cc}^l \right| - 1$$



## Illustrative Numerical Examples



Physical Network: Modified Telecom Italia Network

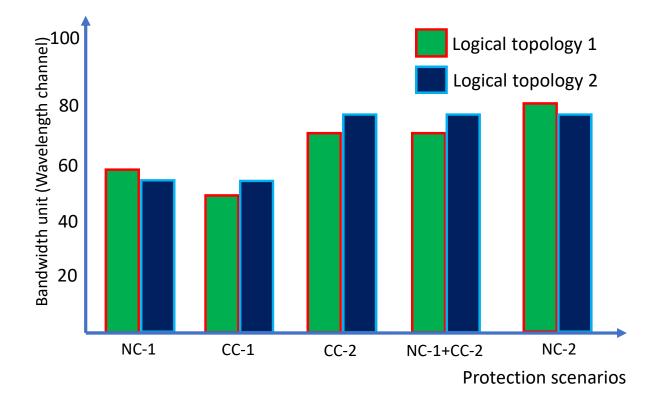


Logical topologies

- ✓ Physical network: 52 nodes, 98 bidirectional links
- ✓ Logical topologies: 7 nodes, 10 bidirectional links (a), and 11 bidirectional links (b)



## Illustrative Numerical Examples



Ongoing: generalizing scenarios

- ✓ NC > CC
- ✓ NC = CC
- ✓ NC not possible but CC possible
- ✓ DC location impact

Our new approach is more generic (arbitrary n) and more scalable than those in [10], [11]



- [10] M. F. Habib, M. Tornatore, and B. Mukherjee, "Fault-tolerant virtual network mapping to provide Content Connectivity in optical networks," *Proceedings of OFC*, Mar. 2013.
- [11] A. Hmaity, F. Musumeci, and M. Tornatore, "Survivable virtual network mapping to provide content connectivity against double-link failures," *Proceedings of DRCN*, Mar. 2016.

## Number of Variables and Constraints

- Variable (only one):  $f_{ij}^{st}$  Capacity Constr.
- $\sum_{st \in E_L} f_{ij}^{st} \le F_{ij} \times W, \forall ij \in E_P$

# variables:  $L_P \times L_L = 196 \times 20 = 3920$ 

# constr.:  $L_P = 196$ 

# constr.:  $N_P \times L_L = 52 \times 20 = 1040$ 

Sum up

# constr. (n = 2):  $|P_n| \times |C_{cc}| = 19,110 \times 30 = 573,300$ 

•  $\sum_{j:ji\in E_P} f_{ji}^{st} - \sum_{j:ij\in E_P} f_{ij}^{st} = \begin{cases} -1 \text{ if } i = s \\ 1 \text{ if } i = t \\ 0 \text{ otherwise} \end{cases}$  $\forall i \in V_P, \forall st \in E_L$ 

CC-n Constr.

Flow Constr.

• 
$$\sum_{\substack{ij \in P_n^k, st \in C_{cc}^l \\ \forall P_n^k \in P_n, \forall C_{cc}^l \in C_{cc}}} f_{ij}^{st} \leq \left| C_{cc}^l \right| - 1$$

IUtali

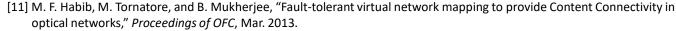
Totally, 3,920 variables, 574,536 constraints (n=2)

## Number of Variables and Constraints Comparison

#### *Note*: Physical topology + Logical topology 1

| Scenarios | Previous works |                 | This work   |  |               |  |
|-----------|----------------|-----------------|-------------|--|---------------|--|
|           | # Variables    | # Constraints   | # Variables |  | # Constraints |  |
| NC-1      | 3,920 [13]     | 25,932 [13]     | 3,920       |  | 25,932        | # var and # constr. reduced by factor of 23 and 12  # var and # constr. reduced by factor of 2 × 10 <sup>3</sup> and 112 |
| CC-1      | 90,220 [11]    | 90,423 [11]     | 3,920       |  | 7,116         |  |
| CC-2      | 8,116,420 [12] | 64,297,083 [12] | 3,920       |  | 574,536       |  |
| NC-1+CC-2 | 8.116,420 [12] | 64,297,083 [12] | 3,920       |  | 599,232       |  |
| NC-2      | NA             | NA              | 3,920       |  | 2,409,096     | Slowly increasing  |

#### # Variables independent of n



<sup>[12]</sup> A. Hmaity, F. Musumeci, and M. Tornatore, "Survivable virtual network mapping to provide content connectivity against double-link failures," *Proceedings of DRCN*, Mar. 2016.

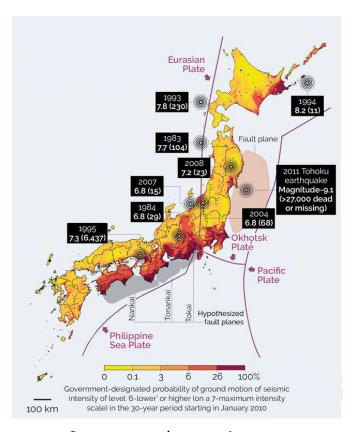
<sup>[13]</sup> E. Modiano and A. Narula-Tam, "Survivable lightpath routing: a new approach to the design of WDM-based networks," *IEEE Journal on Selected Areas in Communications*, vol. 20, no. 4, pp. 800–809, May 2002.



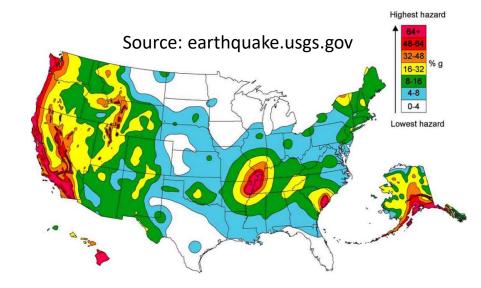
# Chapter 3: Ongoing and Future Research



## Non-Uniform Risk Probability



Source: earthmagazine.org



- ✓ Info from various sources (i.e., geology, climatology, transportation, and environmental science) should be used to determine probability
- ✓ Equipment failure probability due to disasters depends on distance to disaster epicenter, link length, intersection with disaster region



## Flexible Content-Connectivity Protection Plan

#### **Customer:**

- ✓ Require content connectivity for set of nodes (e.g., offices)
- ✓ Demand survivability against large-scale failures

#### Operator:

- ✓ Must satisfy customer's requirements
- ✓ Question: Fixed logical topology?

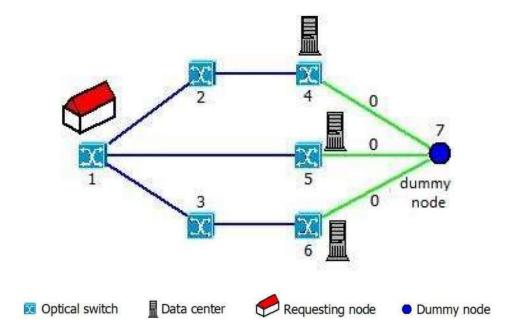
#### Flexible Content Connectivity Protection:

- ✓ Optimally design (lowest cost) logical topology with options
- ✓ Add more datacenters
- ✓ Add more logical links



## Dummy Node Approach for Content Connectivity

- CC-n: Find n+1 link-disjoint paths from content-requesting node (node 1) to dummy node through datacenters
- We expect fast optimal solutions (higher scalability)

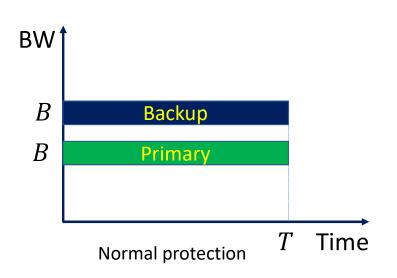


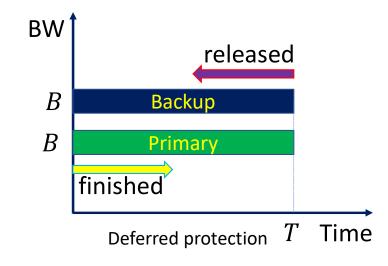


## **Deferred Protection for Content Connectivity**

Major delay of a large file transmission:  $\frac{Packet\ size}{Link\ BW} = \frac{L}{B} = T$  (s)

Do we need content connectivity protection for entire T?









## Acknowledgement

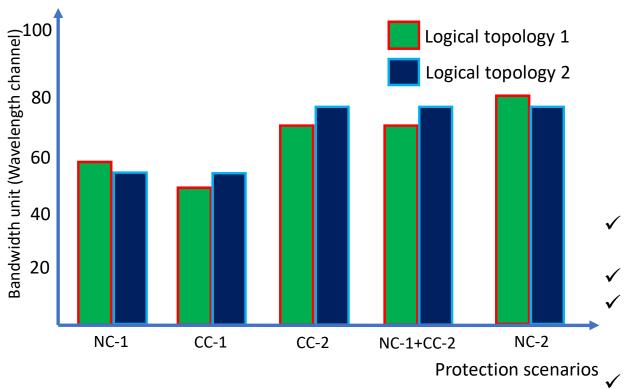
• NSF-JUNO2 CNS-1818972

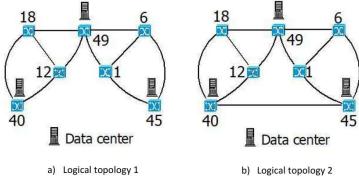


# **Backup Slides**



## Why NC-1 (Topo 1) > NC-1 (Topo 2)?



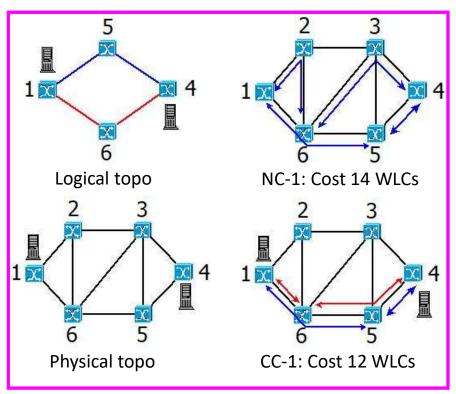


Logical topologies

- √ 40-45 link makes topology hard to be disconnected
- ✓ 6 wavelength channels for 40-45
  - But save 8 wavelength channels for other links taking shorter paths (less strict conditions)
- ✓ NC-1 (Topo 2): 2 wavelength channels less

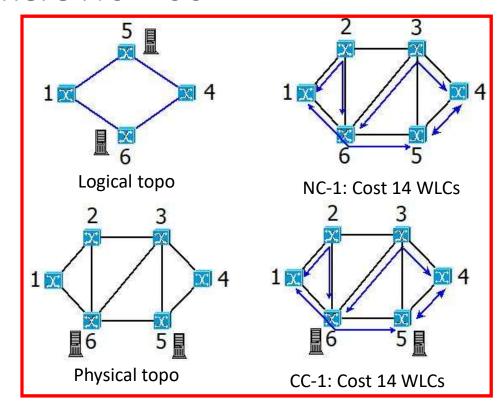


## Scenarios Where NC = CC





- ✓ DCs fanned out
- ✓ Overlapping (i.e., link 5-6)



- ✓ CC = NC
- ✓ DCs inner part

